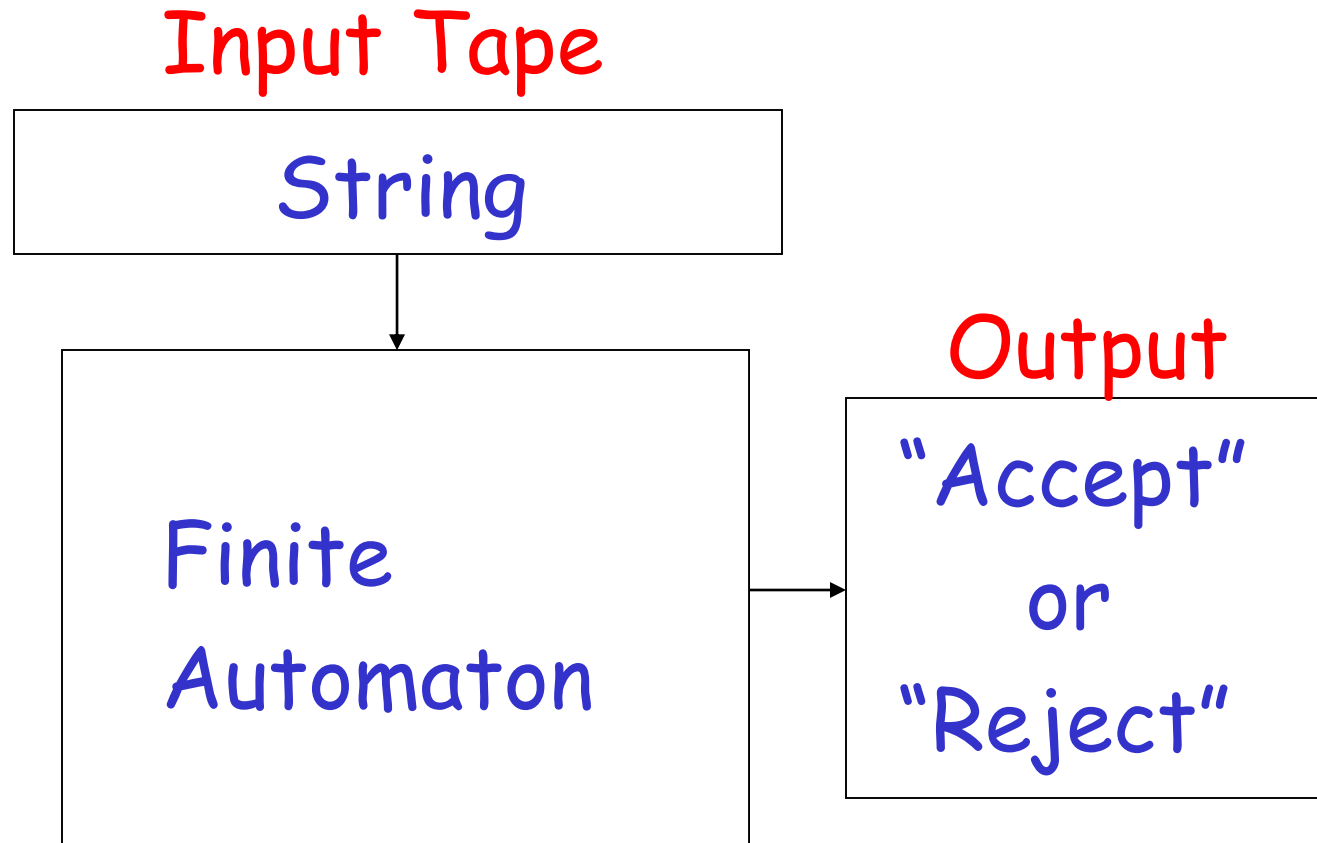


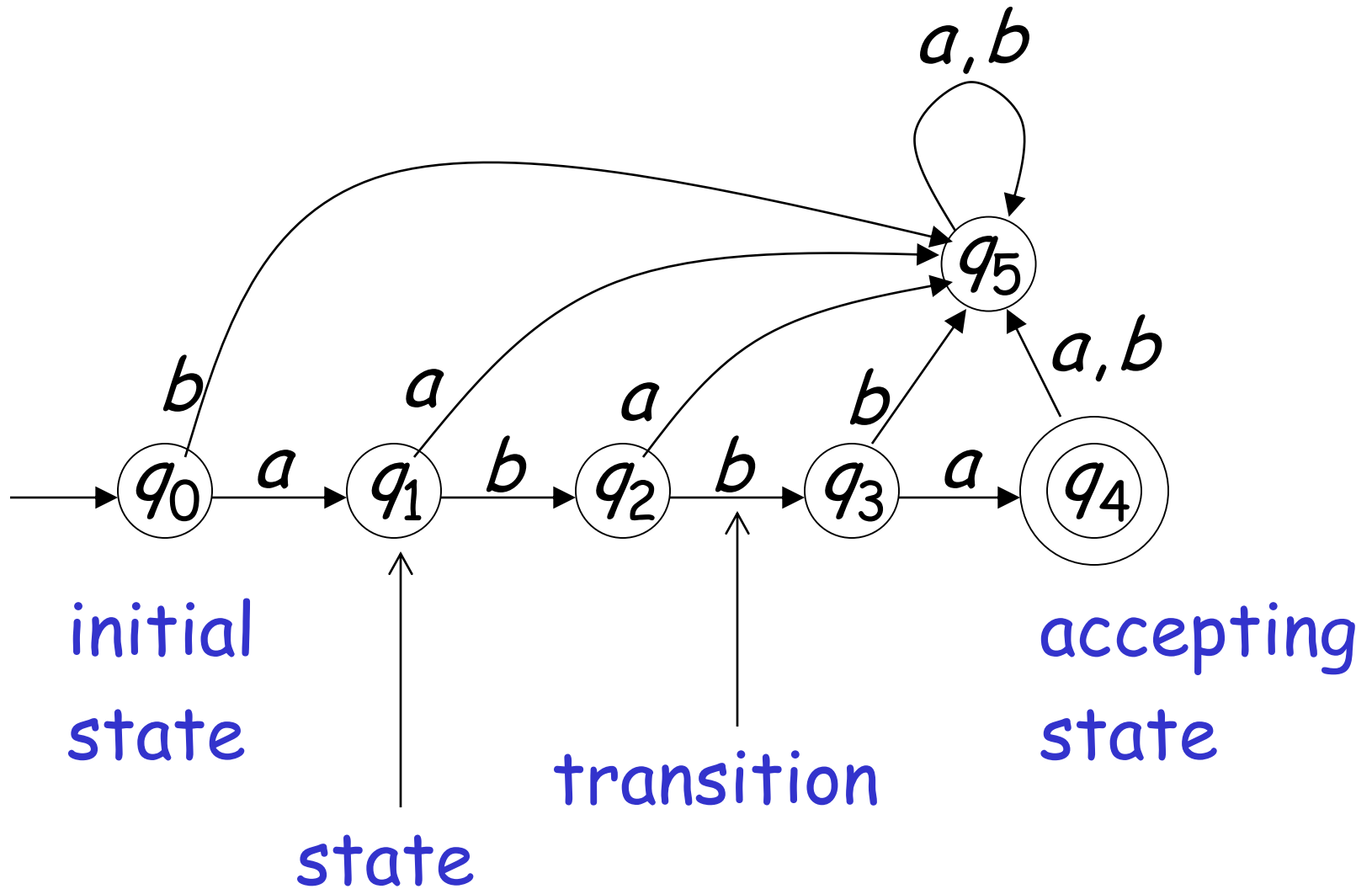
Deterministic
Finite Automata

And Regular Languages

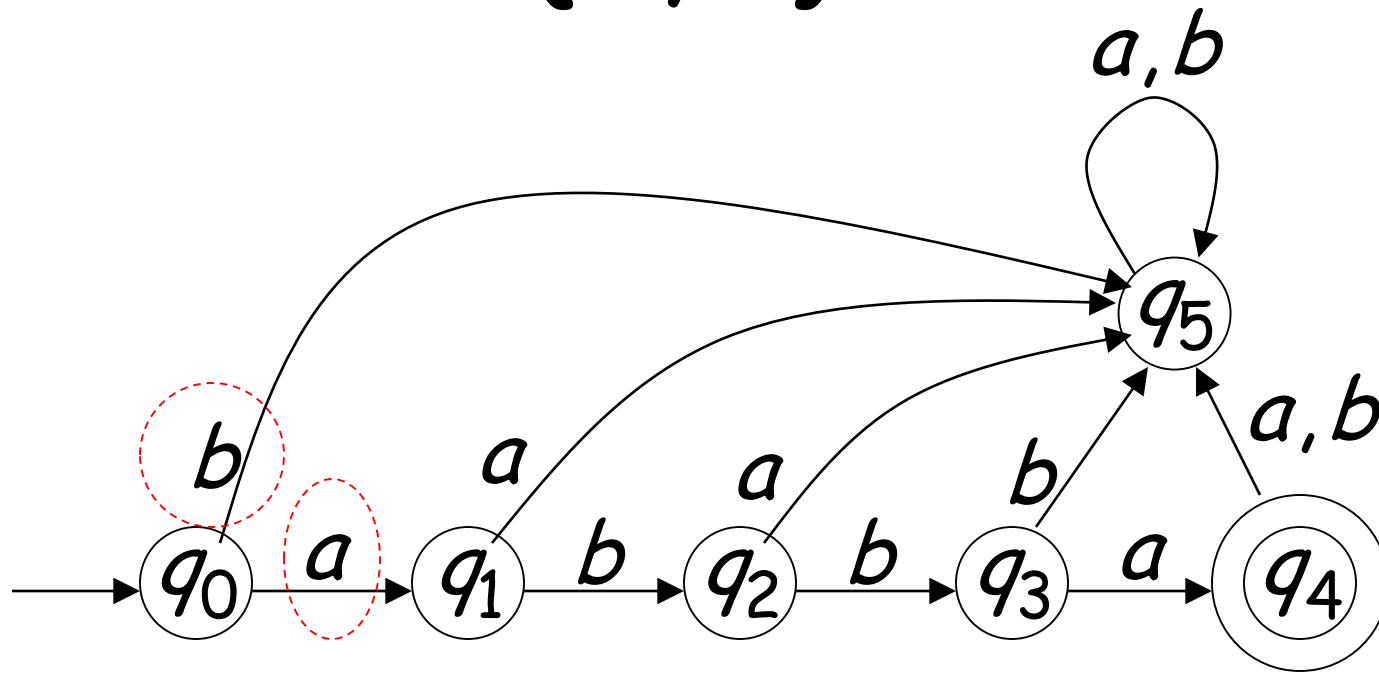
Deterministic Finite Automaton (DFA)



Transition Graph



Alphabet $\Sigma = \{a, b\}$



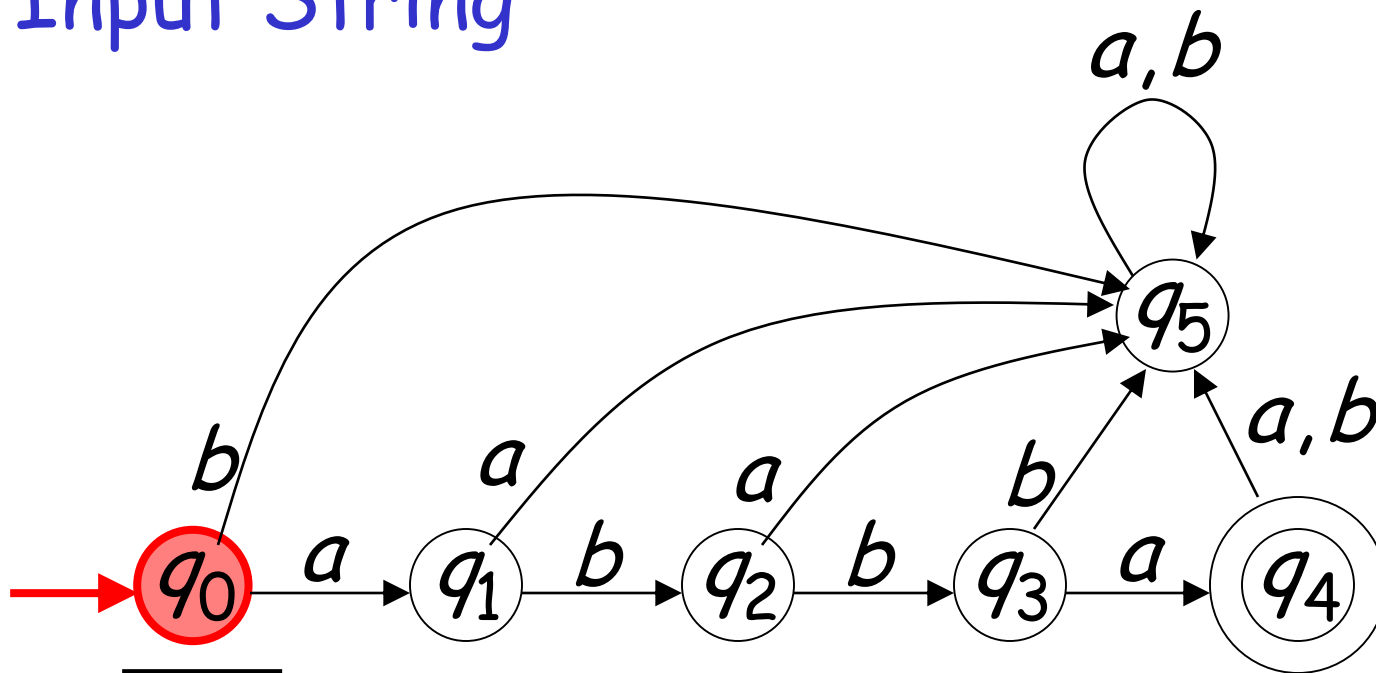
For every state, there is a transition for every symbol in the alphabet

Initial Configuration

head

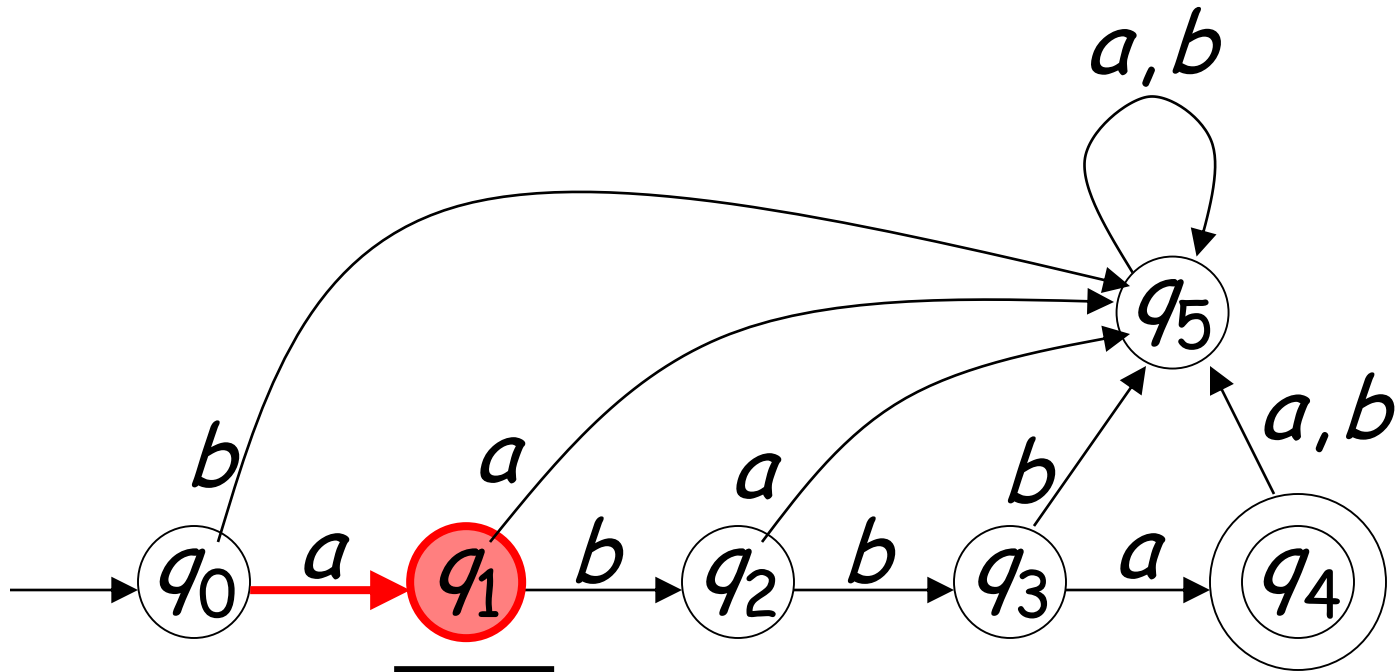
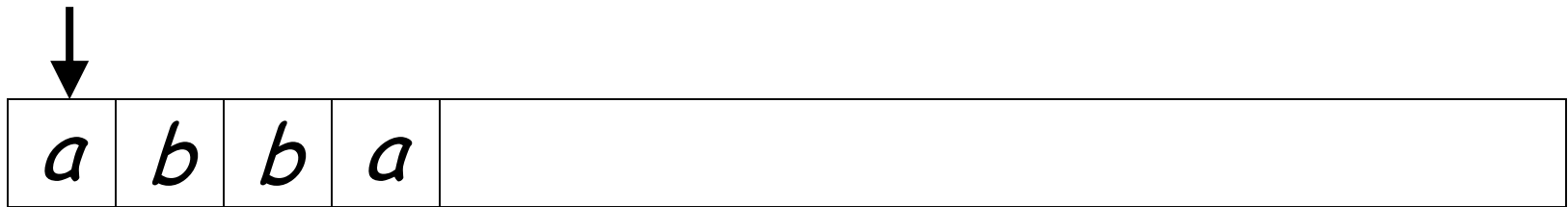


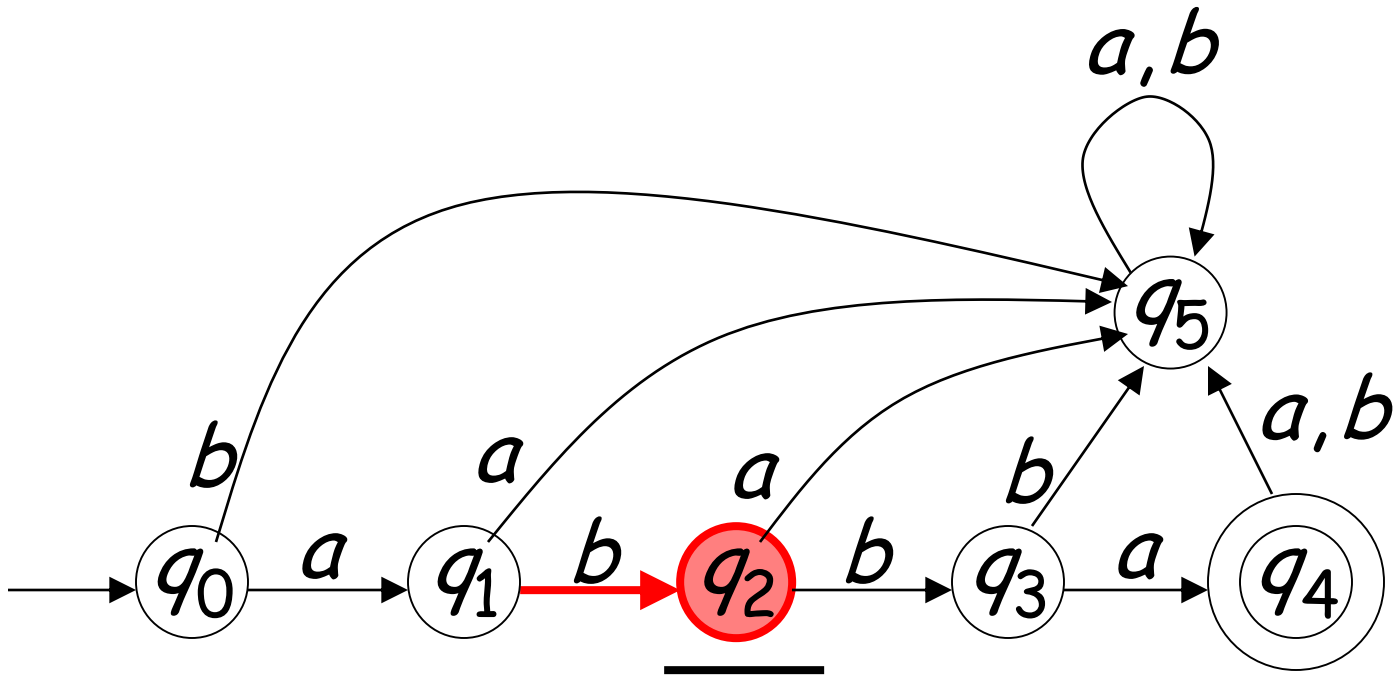
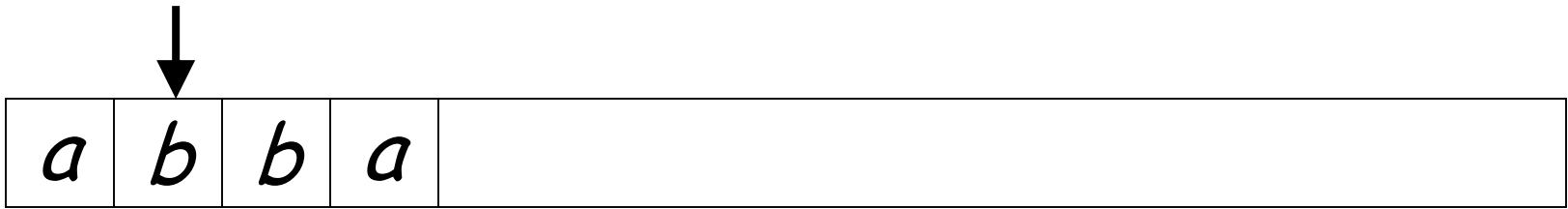
Input String

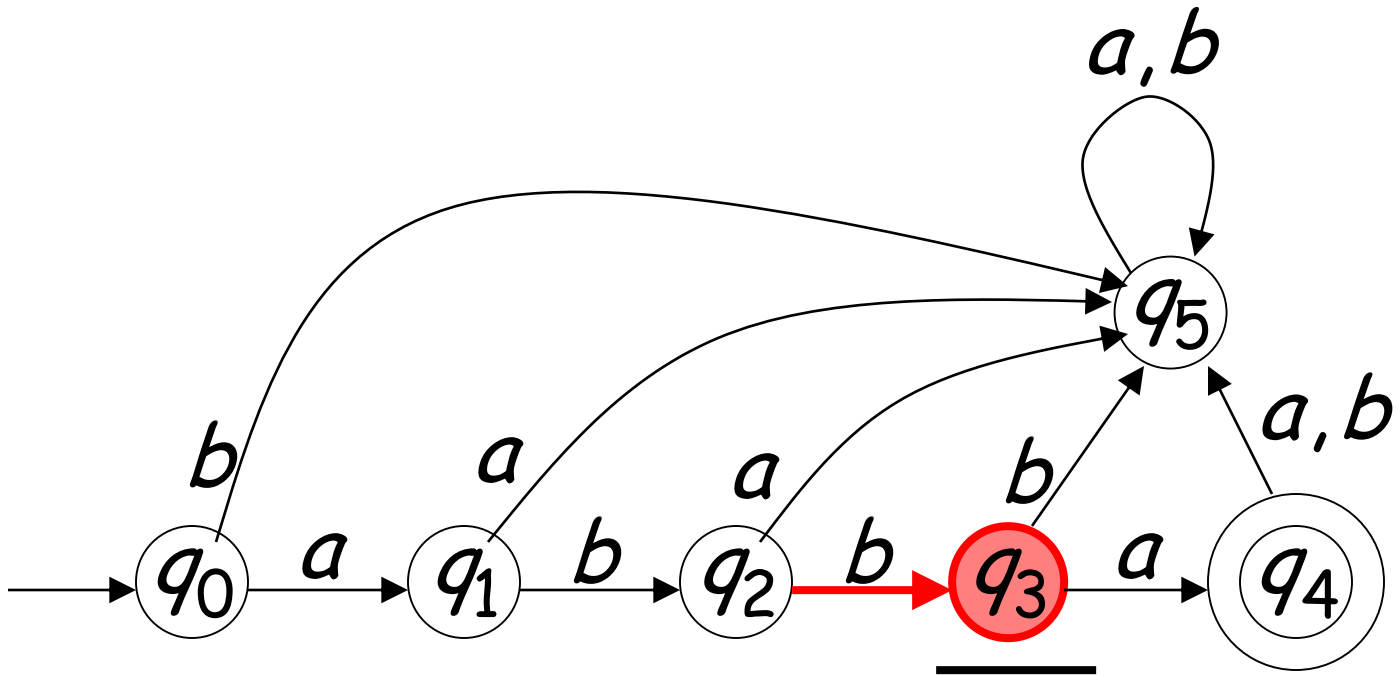
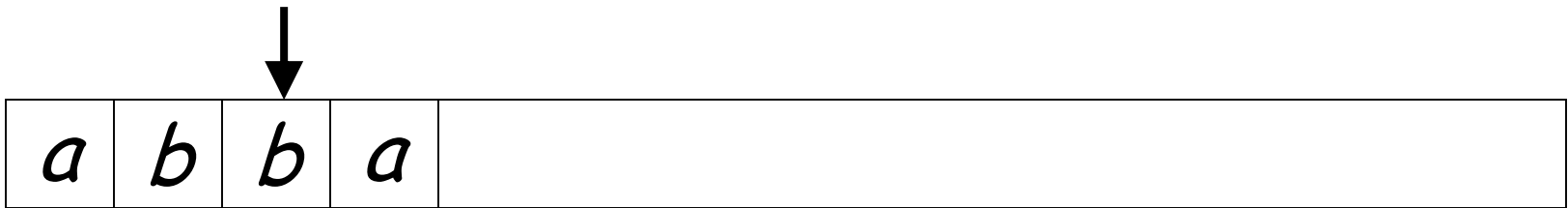


Initial state

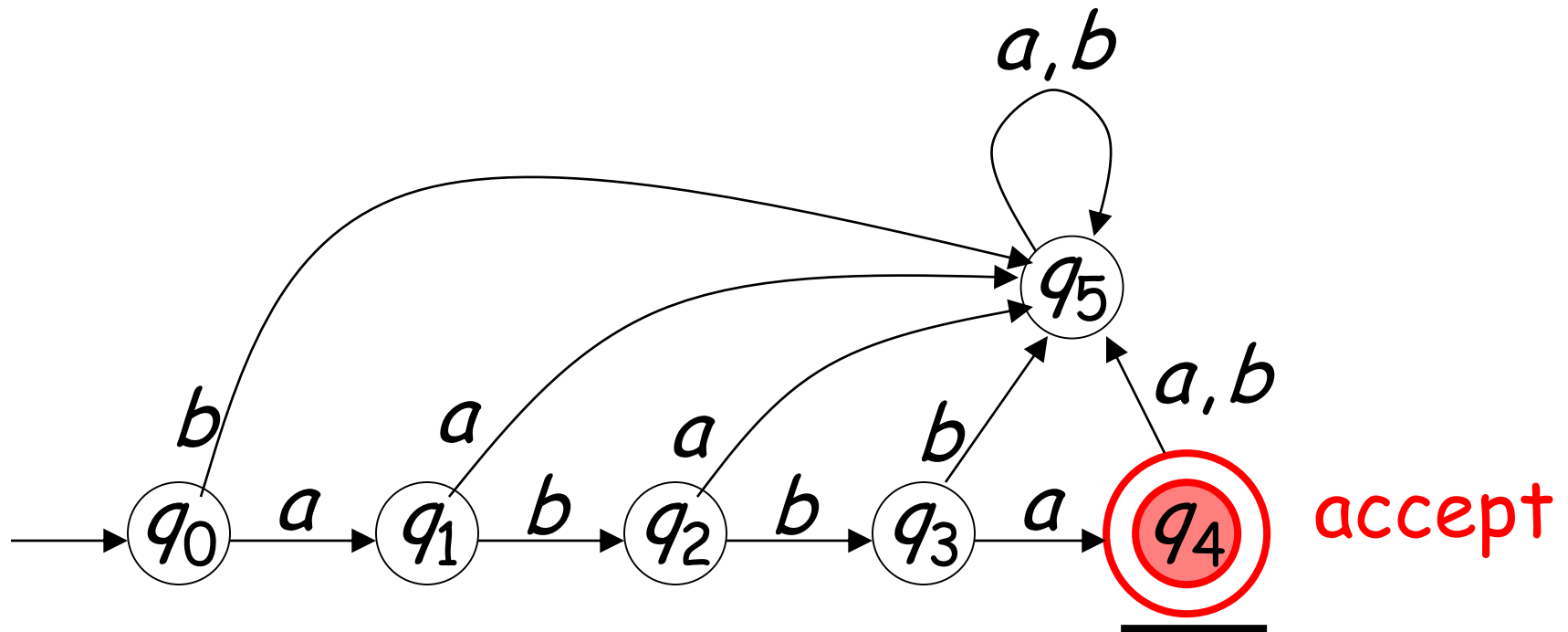
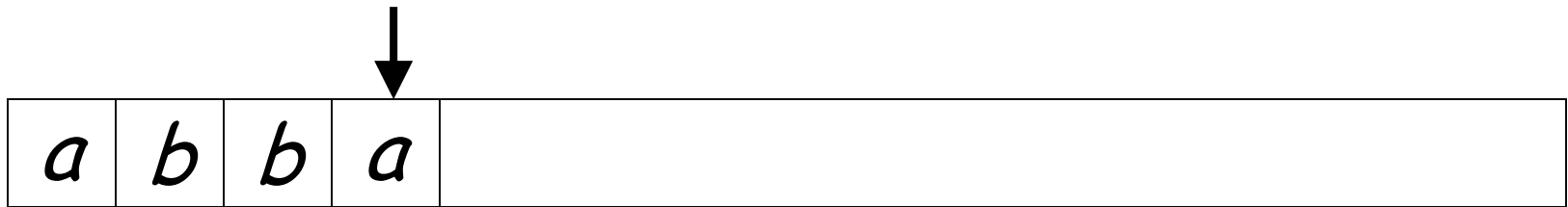
Scanning the Input





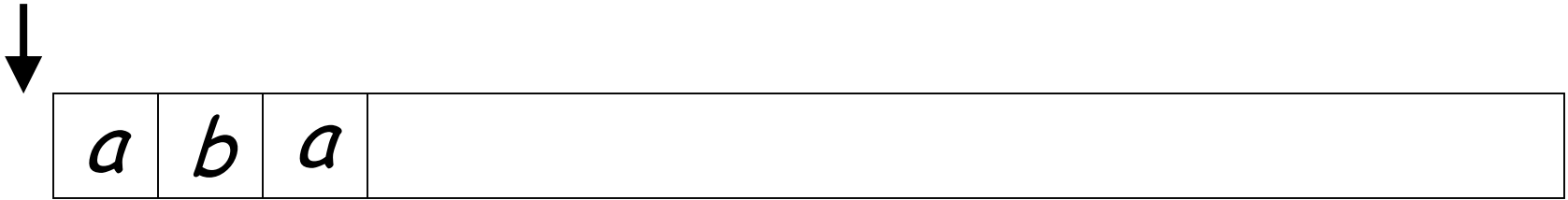


Input finished

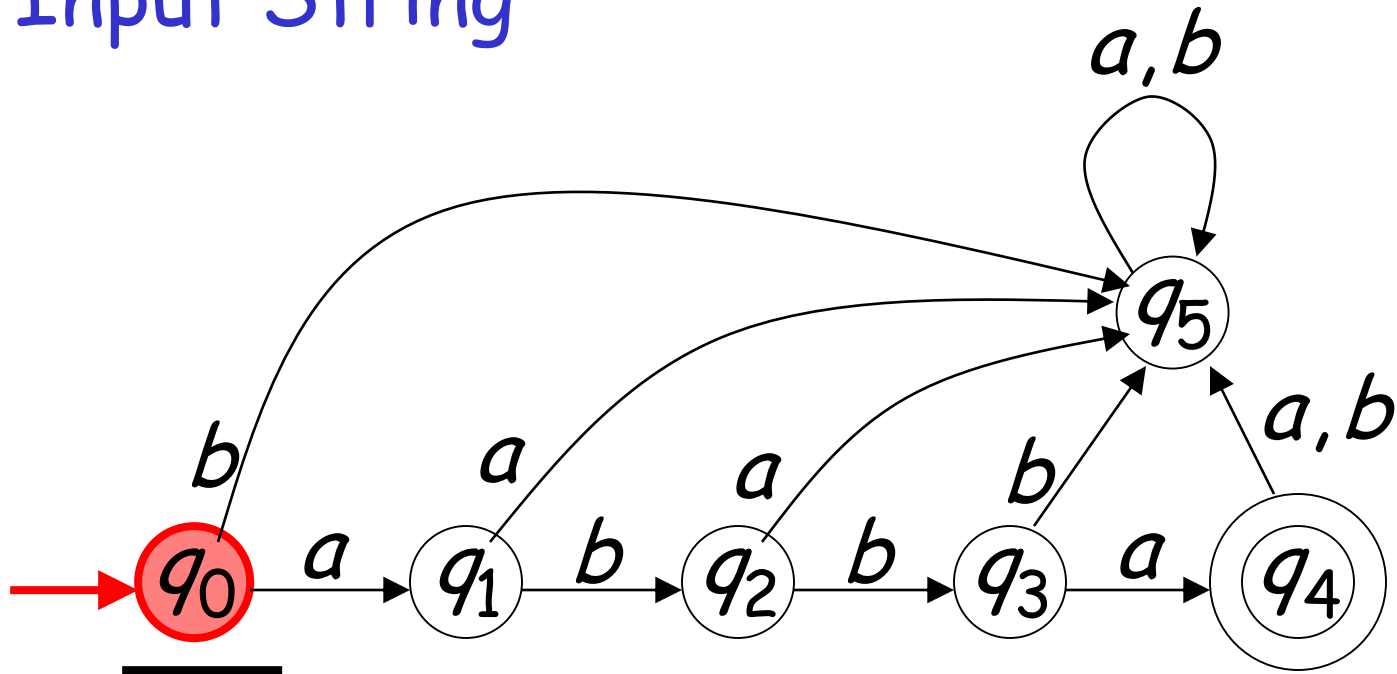


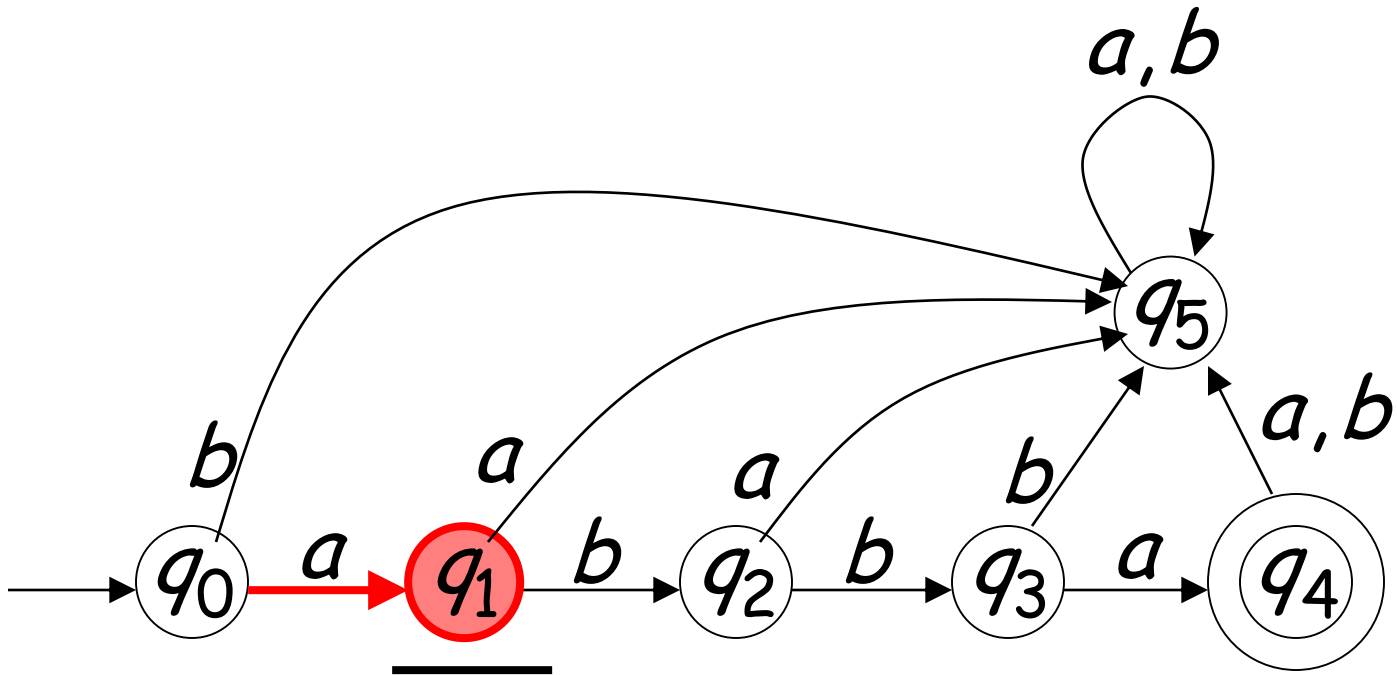
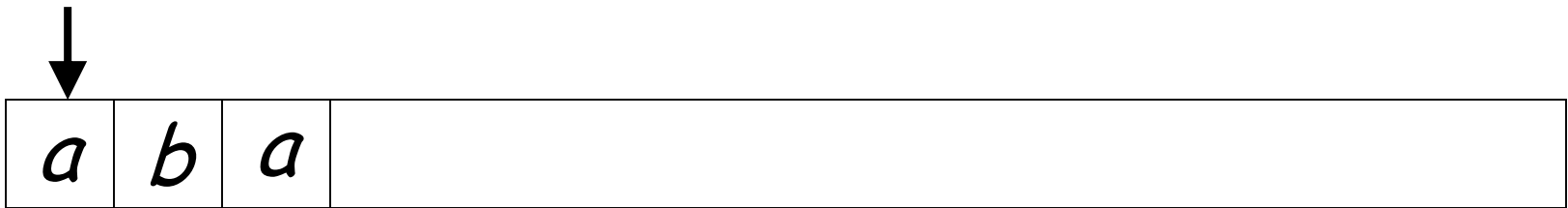
Last state determines the outcome

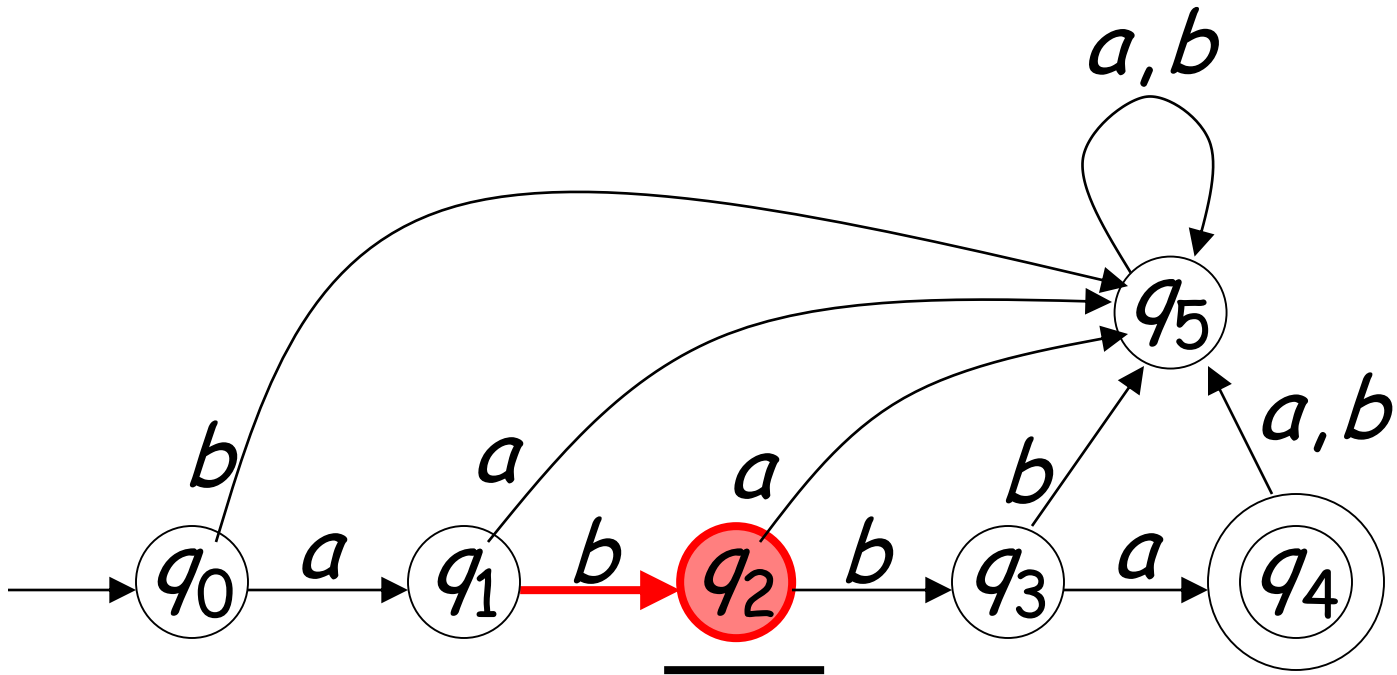
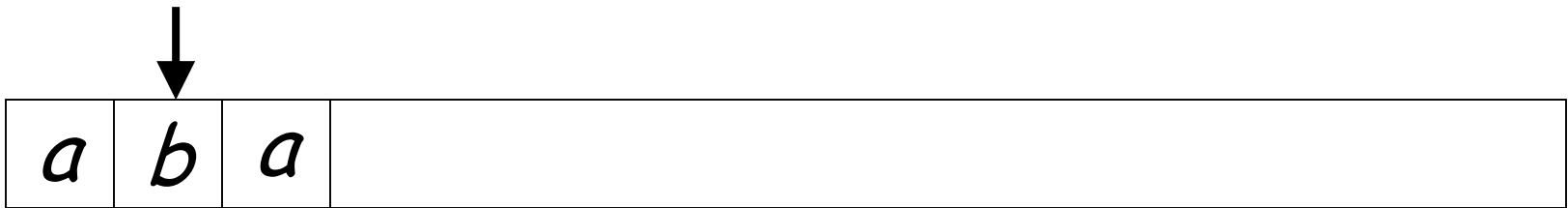
A Rejection Case



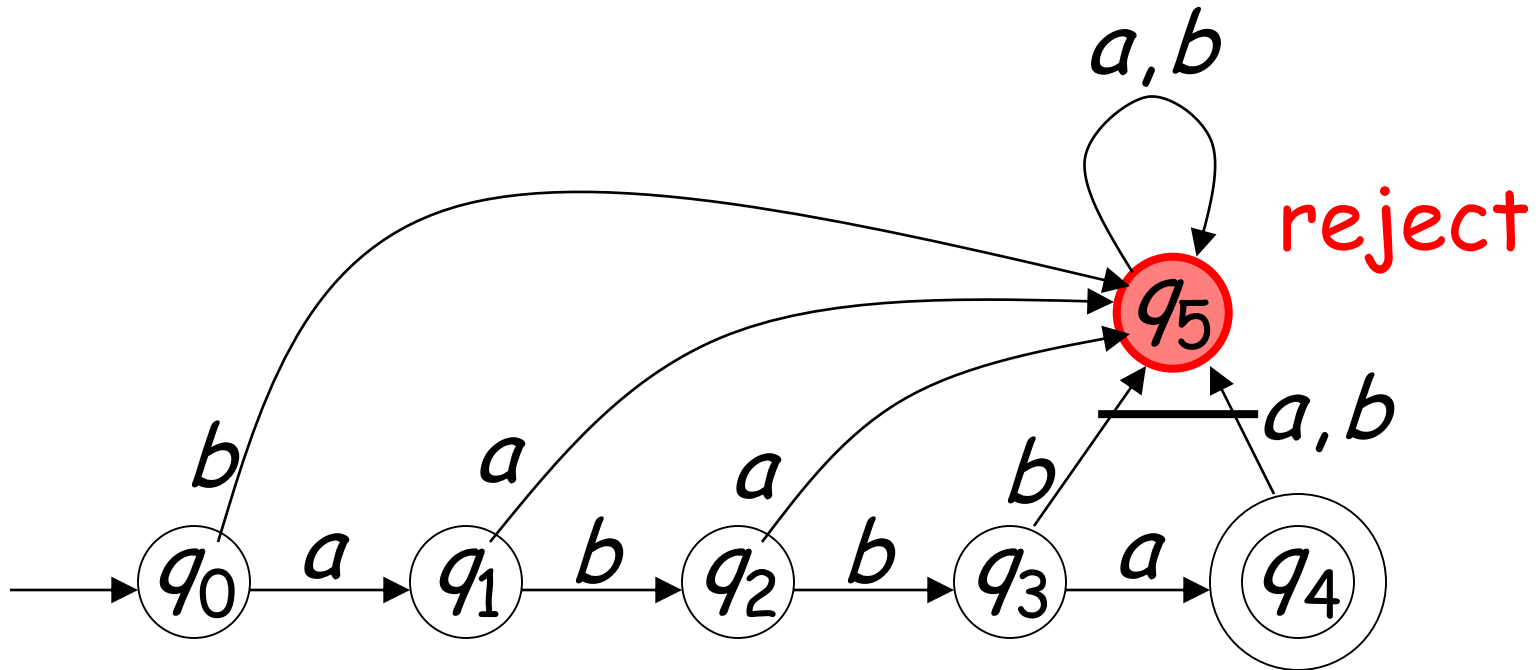
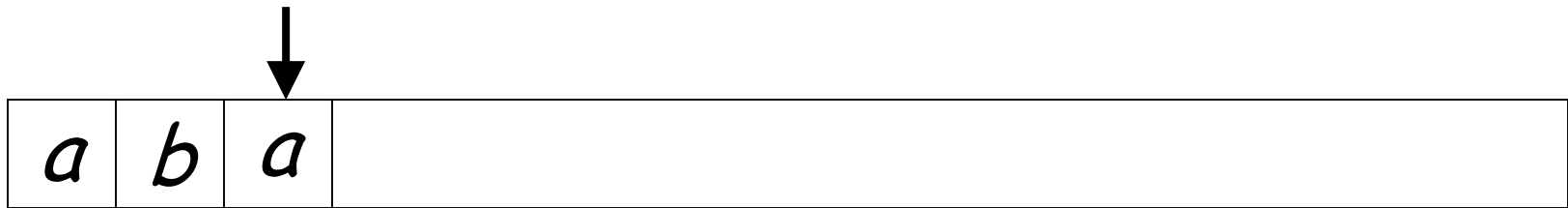
Input String





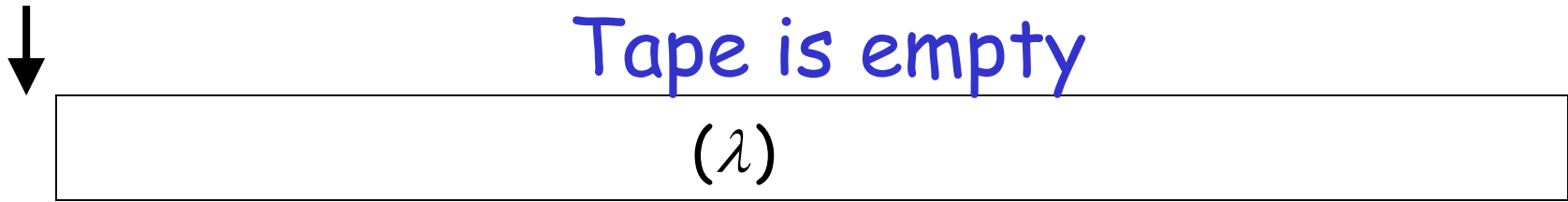


Input finished

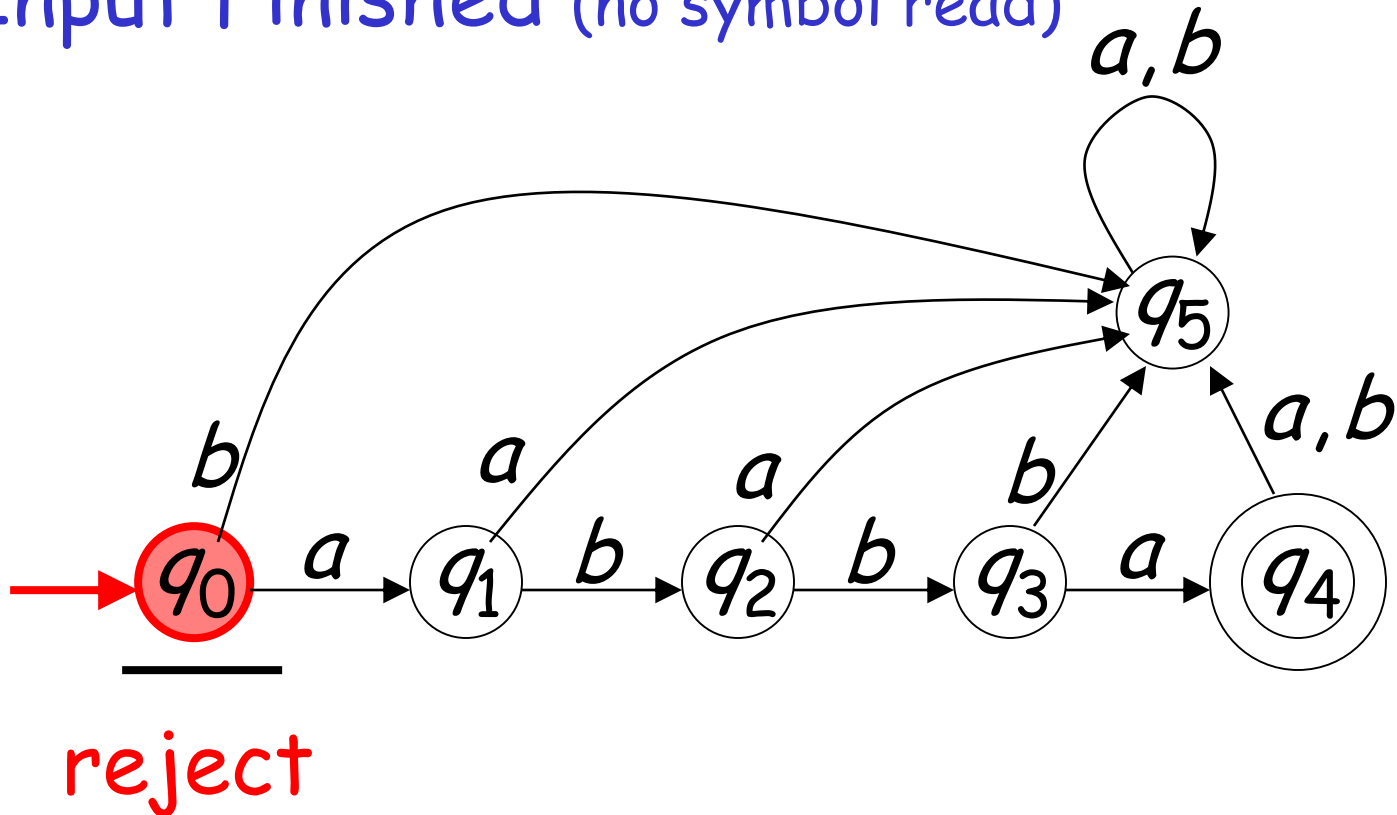


Last state determines the outcome

Another Rejection Case

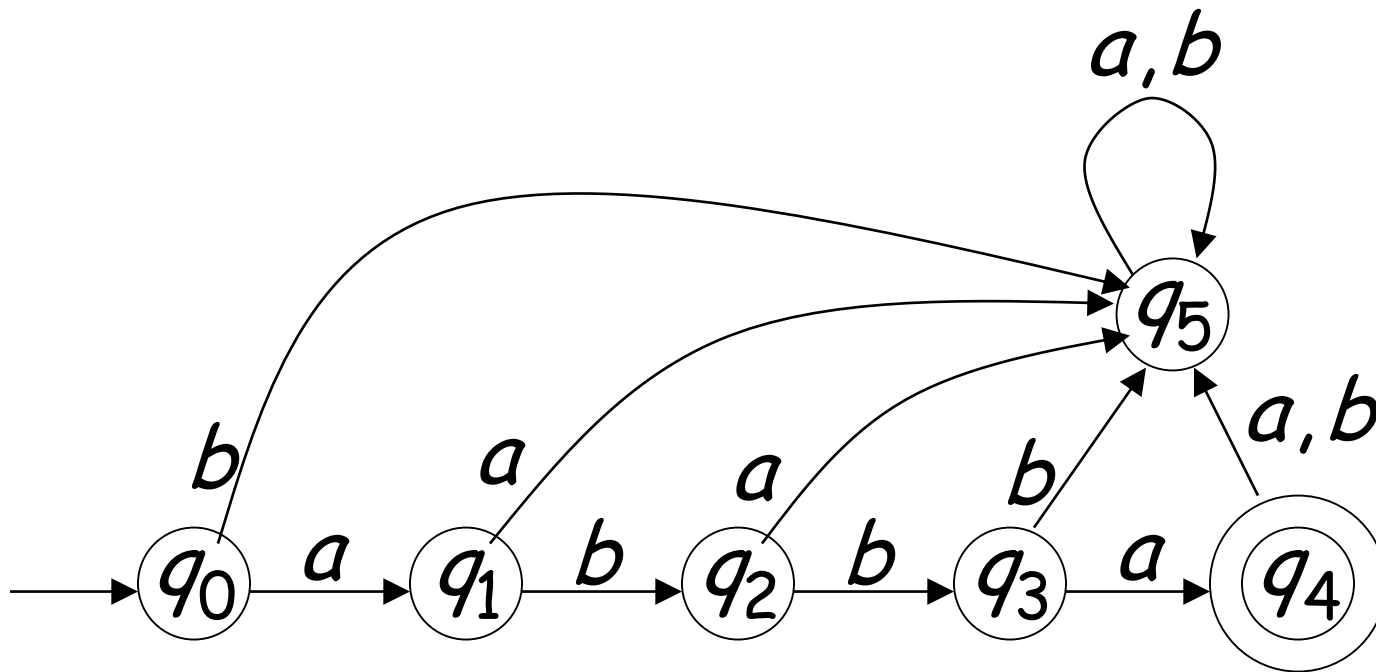


Input Finished (no symbol read)



This automaton accepts only one string

Language Accepted: $L = \{abba\}$



To accept a string:

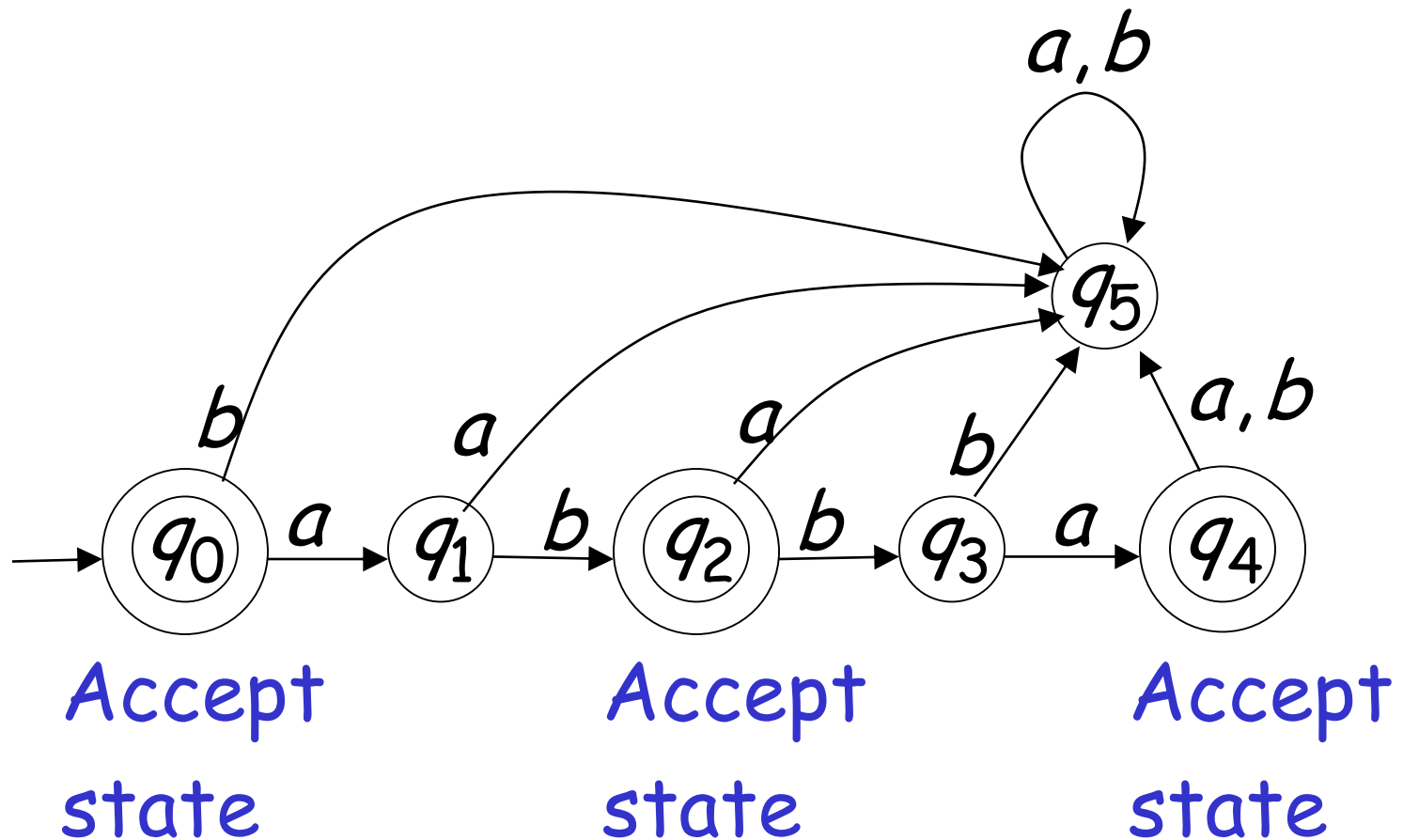
all the input string is scanned
and the last state is accepting

To reject a string:

all the input string is scanned
and the last state is non-accepting

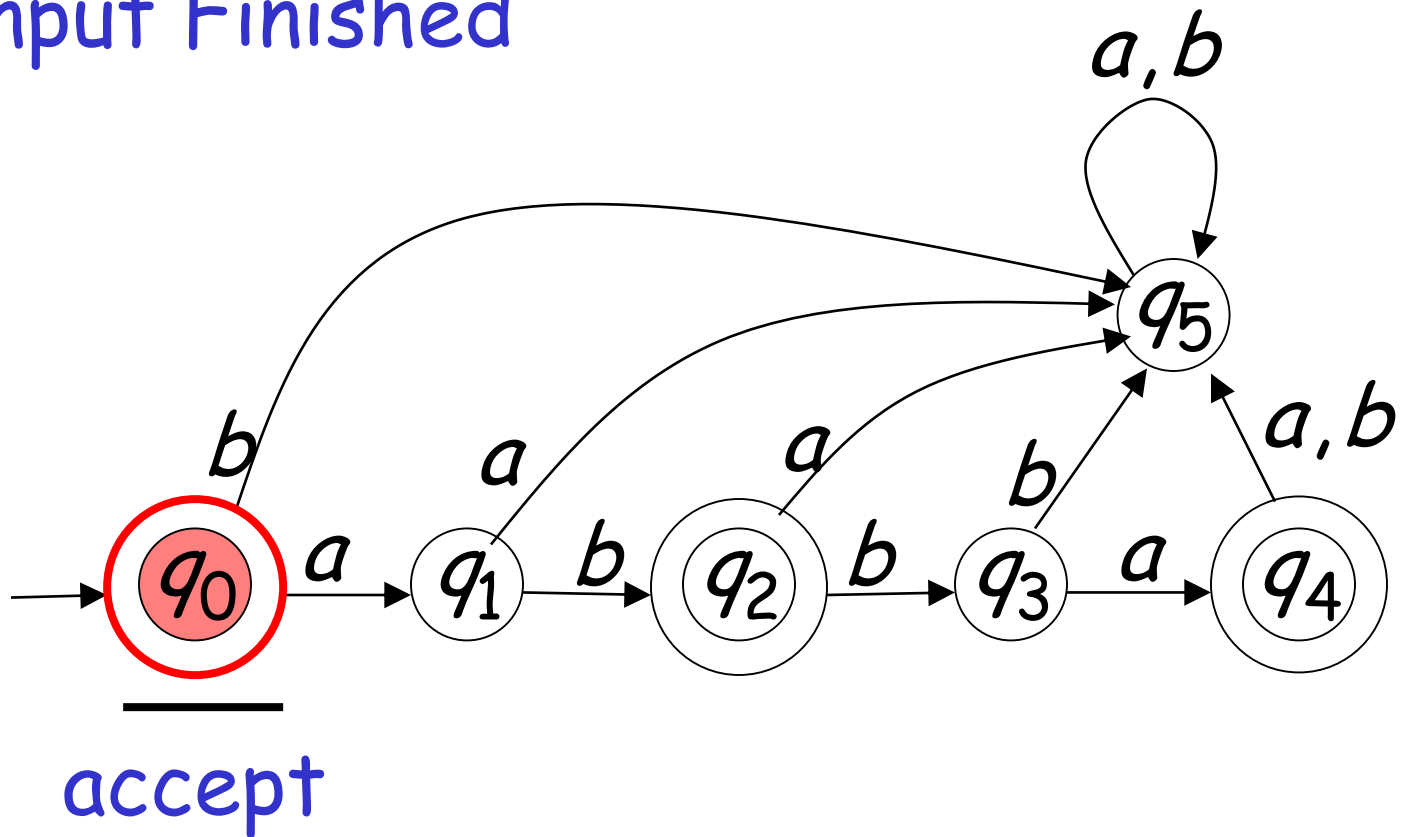
Another Example

$$L = \{ \lambda, ab, abba \}$$

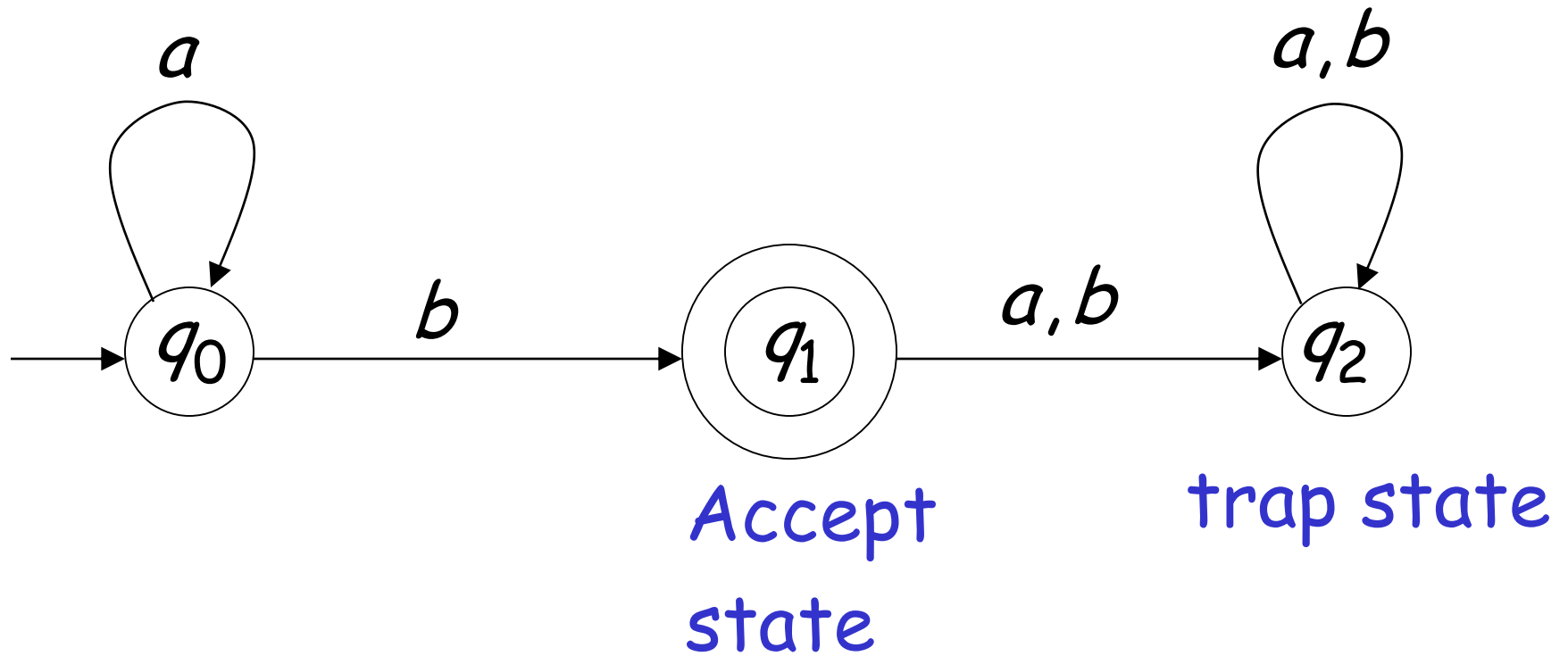


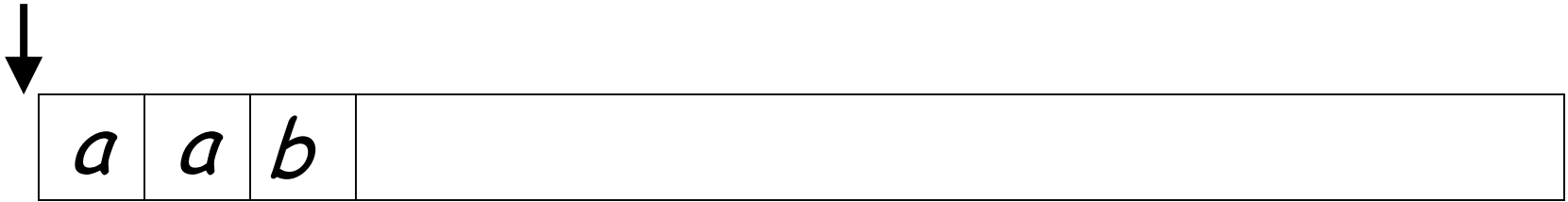


Input Finished

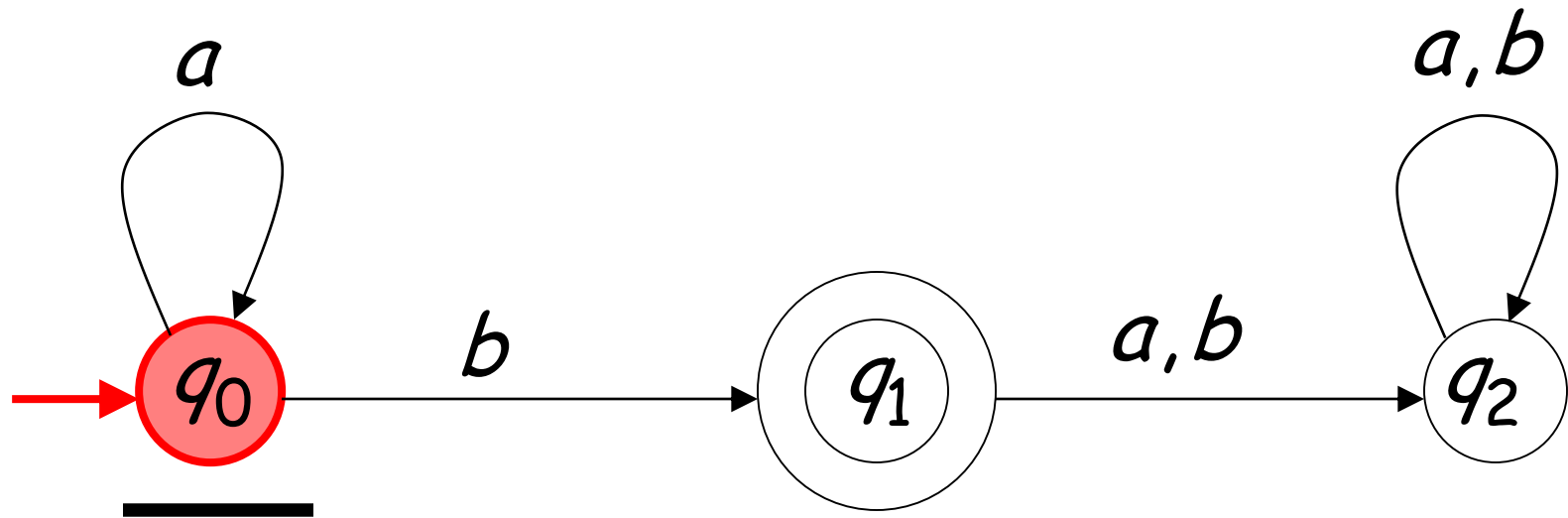


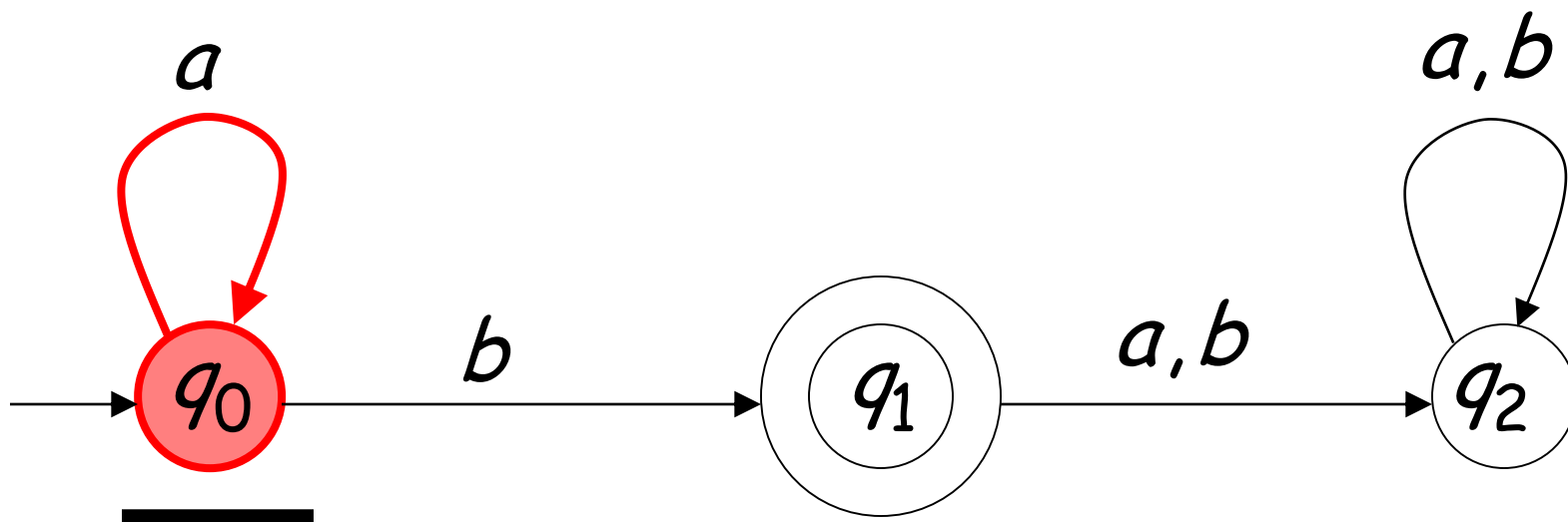
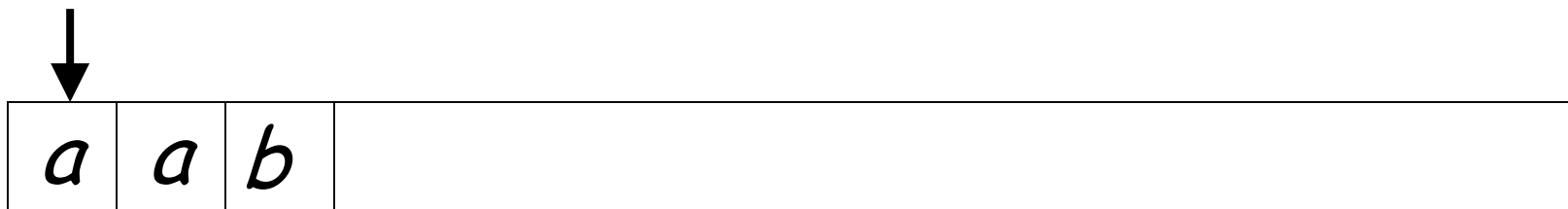
Another Example

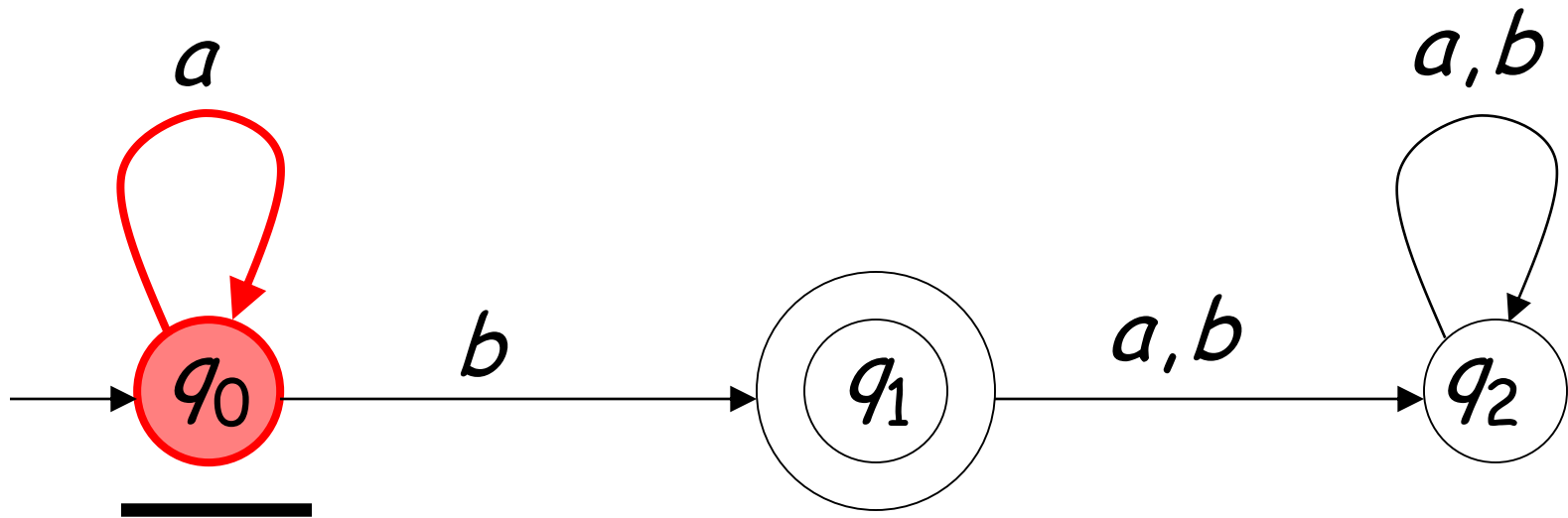
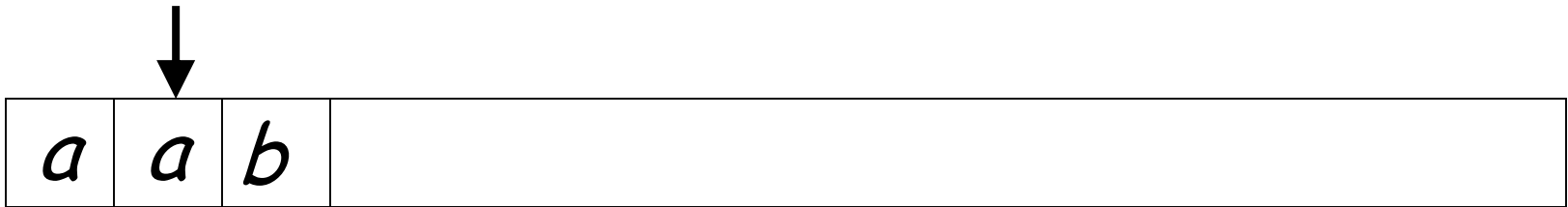




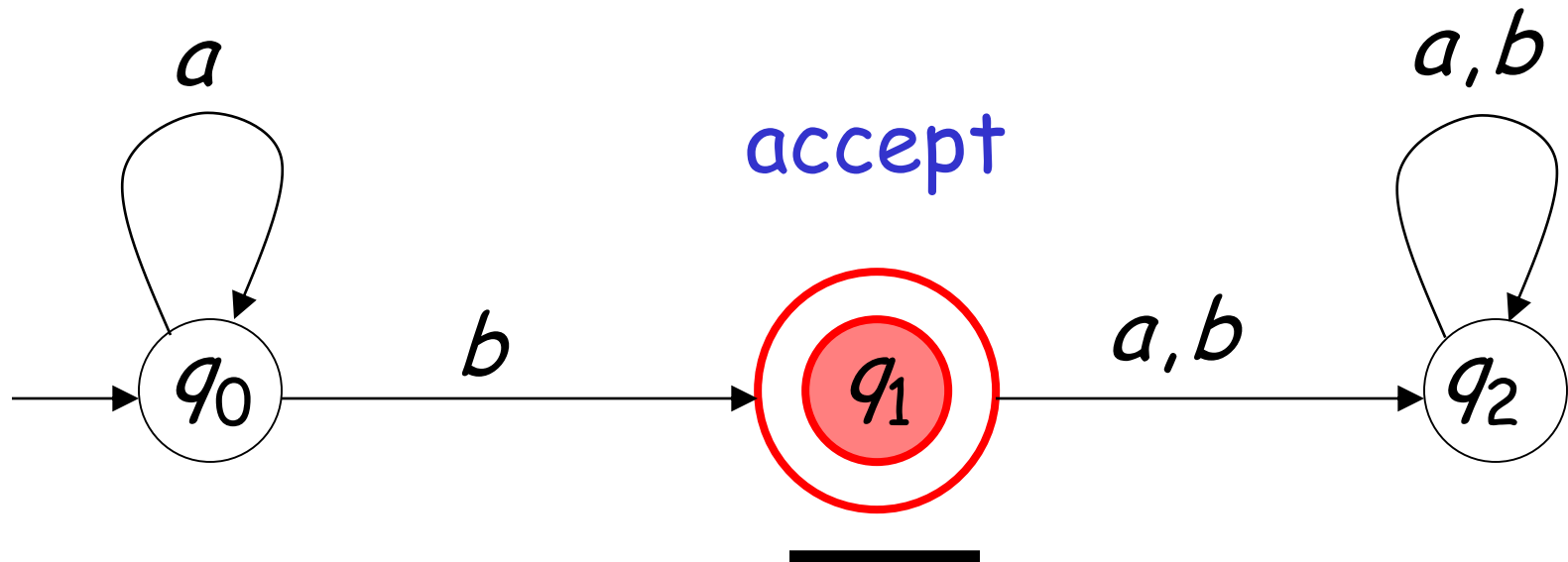
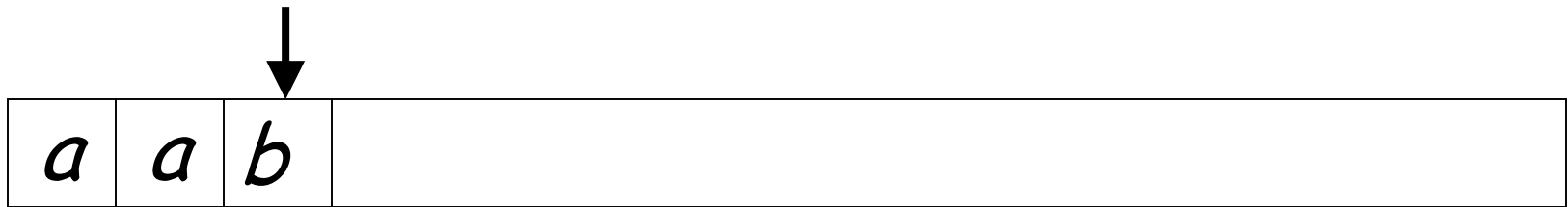
Input String







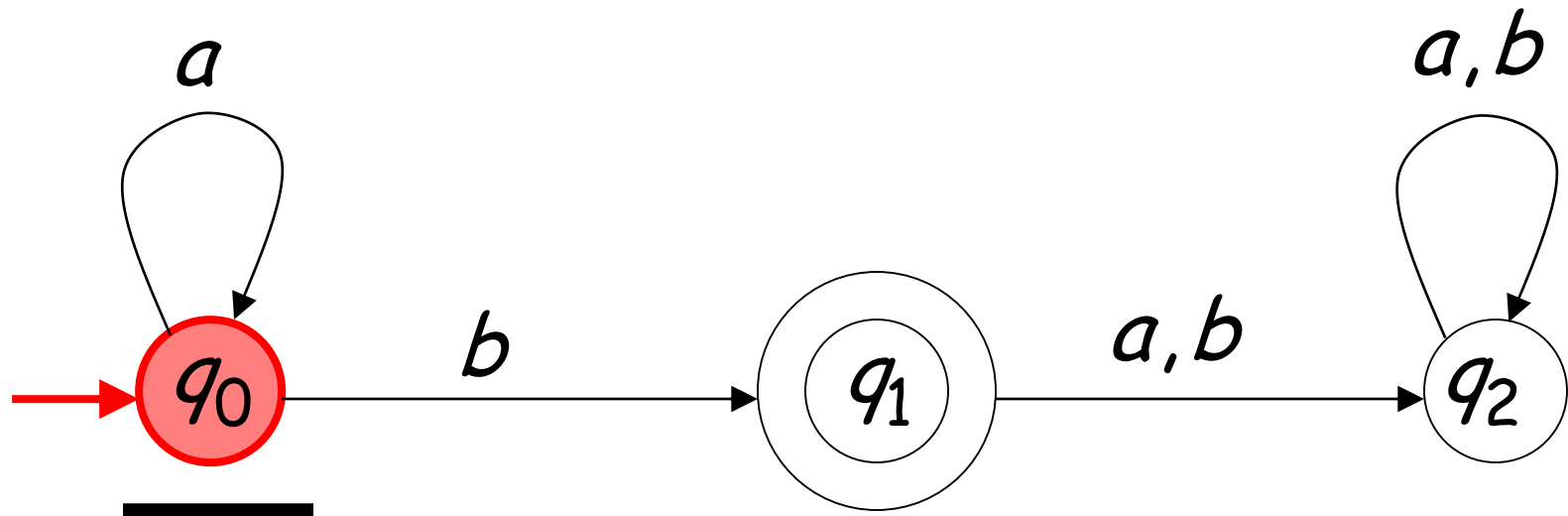
Input finished

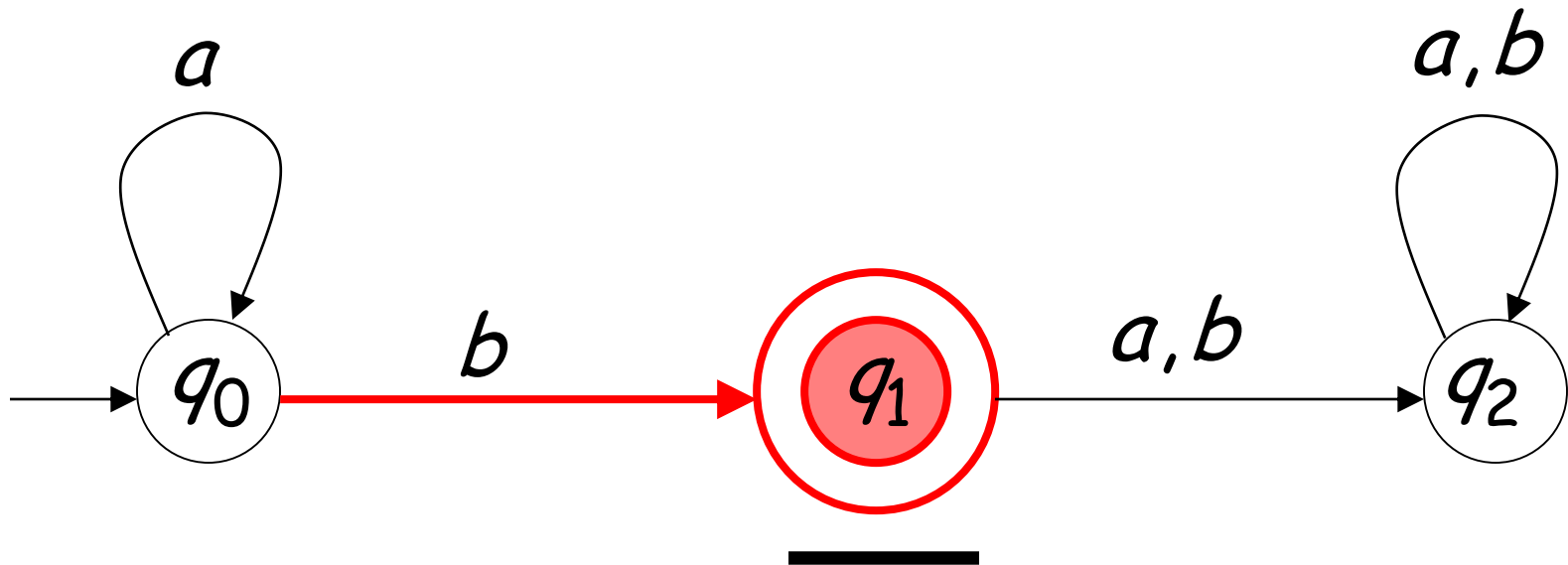
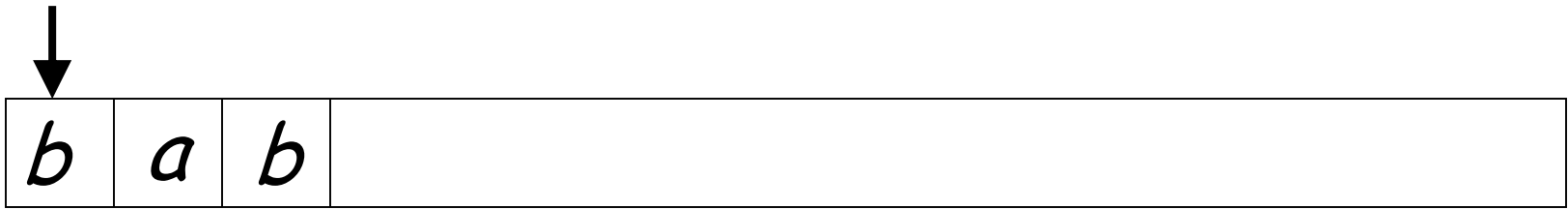


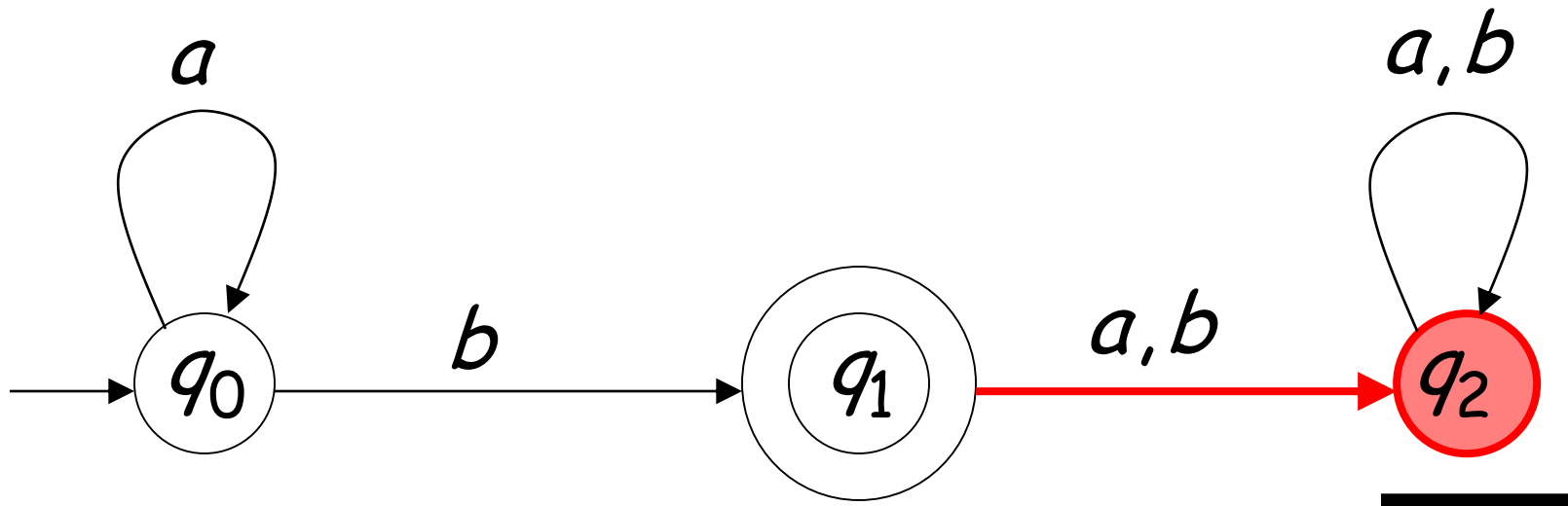
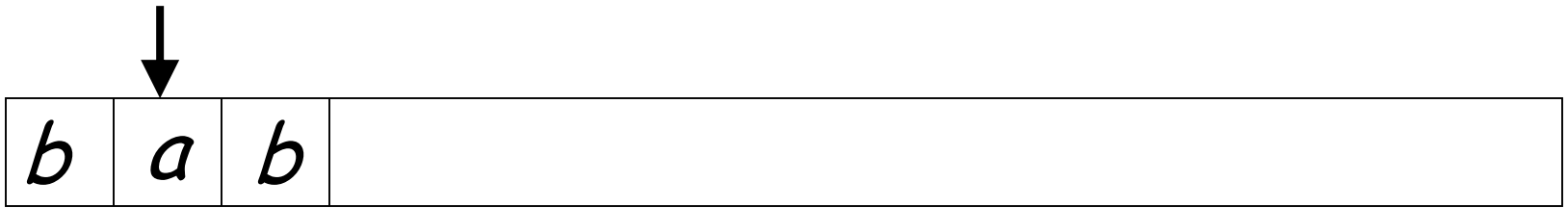
A rejection case



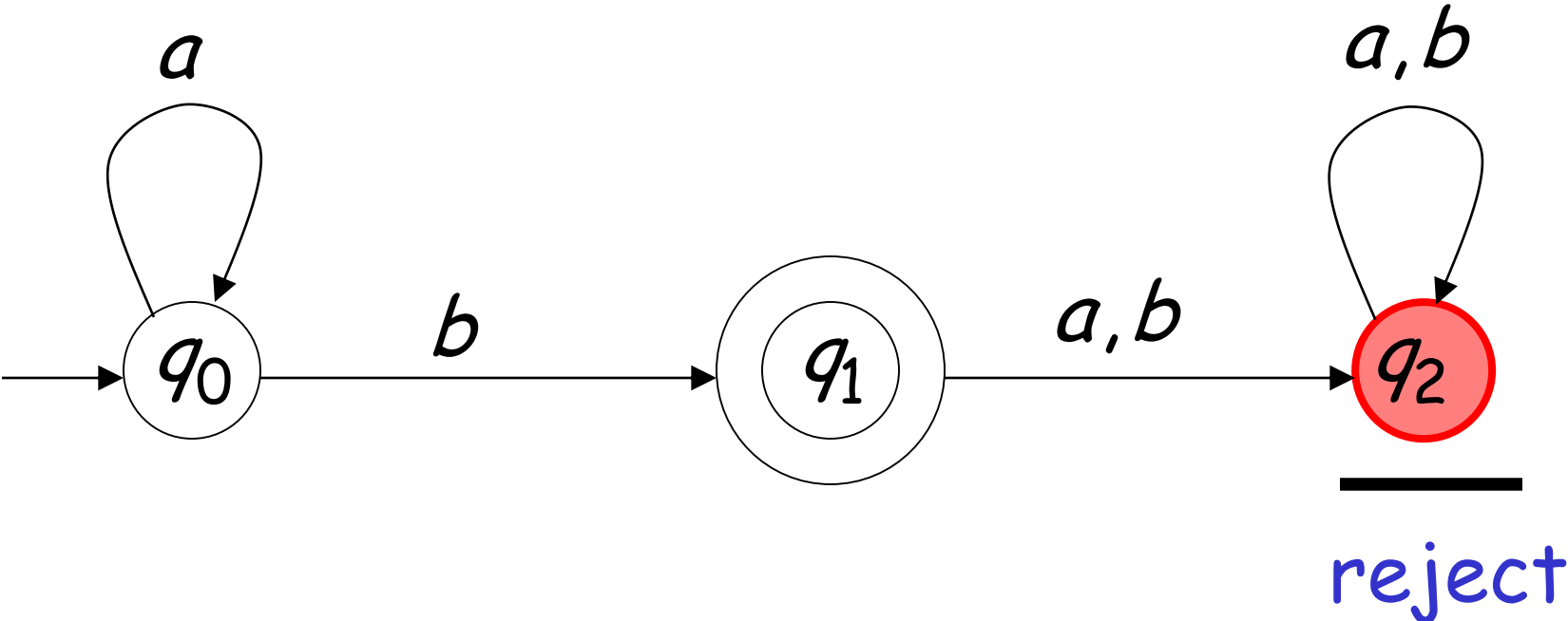
Input String



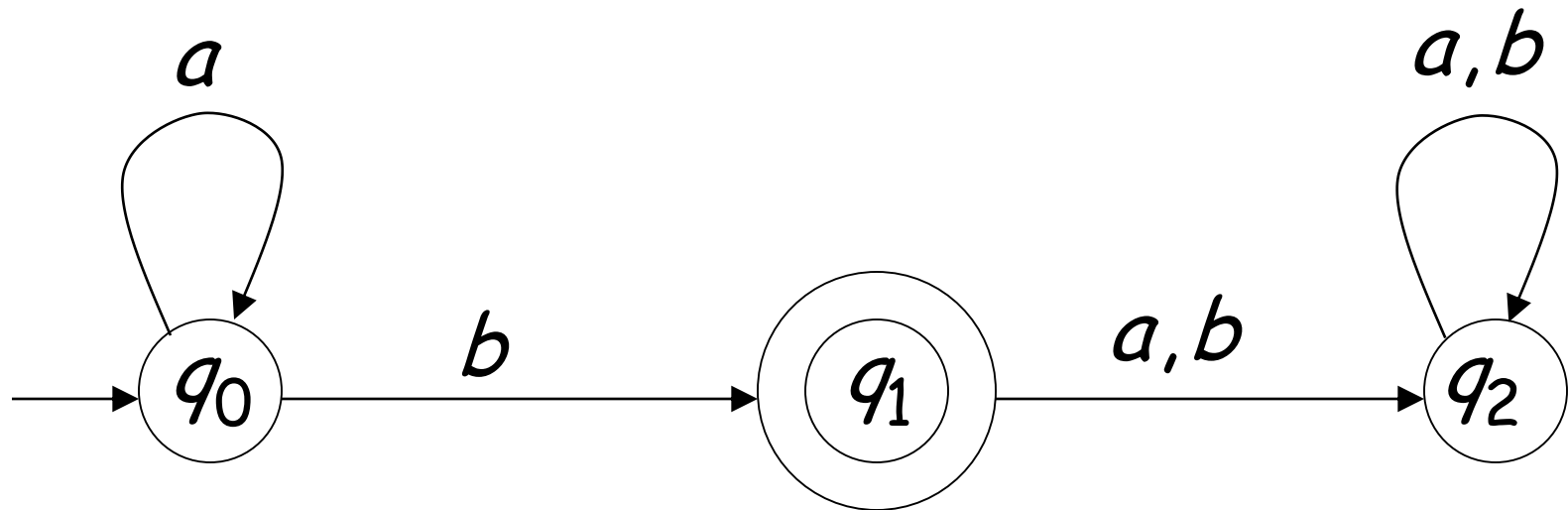




Input finished

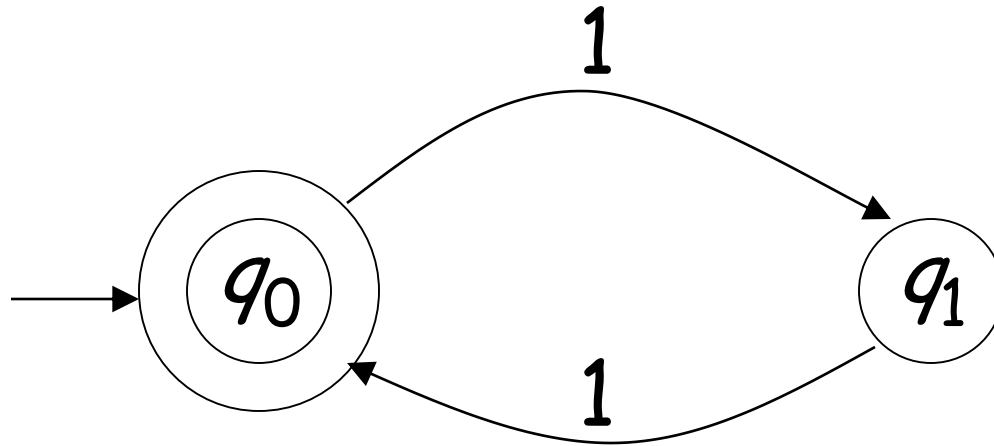


Language Accepted: $L = \{a^n b : n \geq 0\}$



Another Example

Alphabet: $\Sigma = \{1\}$



Language Accepted:

$$\begin{aligned} \text{EVEN} &= \{x : x \in \Sigma^* \text{ and } x \text{ is even}\} \\ &= \{\lambda, 11, 1111, 111111, \dots\} \end{aligned}$$

Formal Definition

Deterministic Finite Automaton (DFA)

$$M = (Q, \Sigma, \delta, q_0, F)$$

Q : set of states

Σ : input alphabet $\lambda \notin \Sigma$

δ : transition function

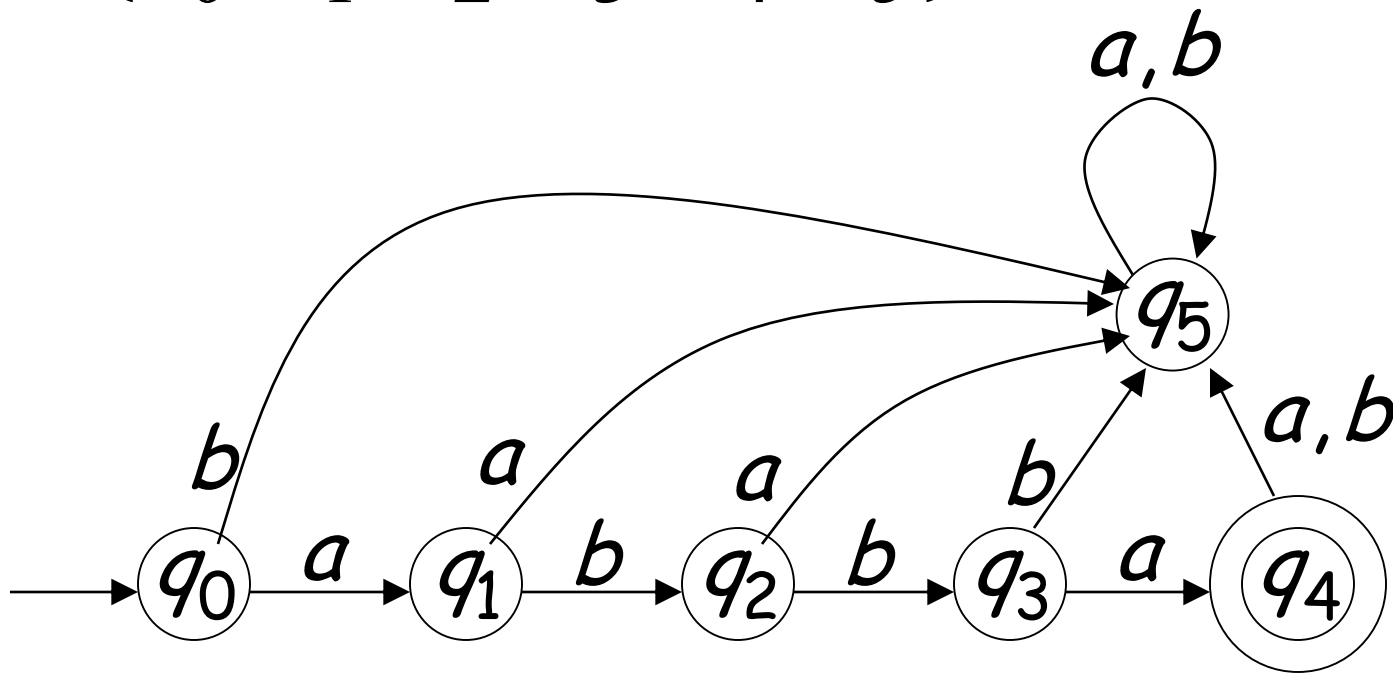
q_0 : initial state

F : set of accepting states

Set of States Q

Example

$$Q = \{q_0, q_1, q_2, q_3, q_4, q_5\}$$

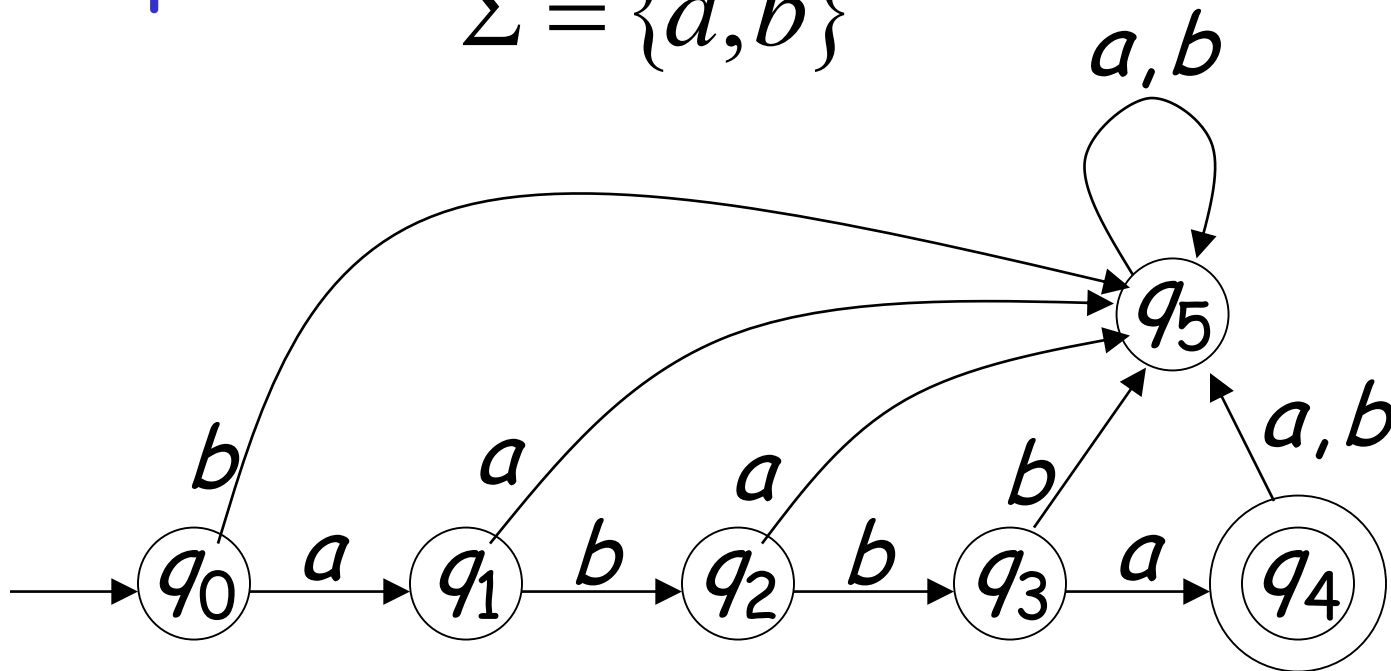


Input Alphabet Σ

$\lambda \notin \Sigma$: the input alphabet never contains λ

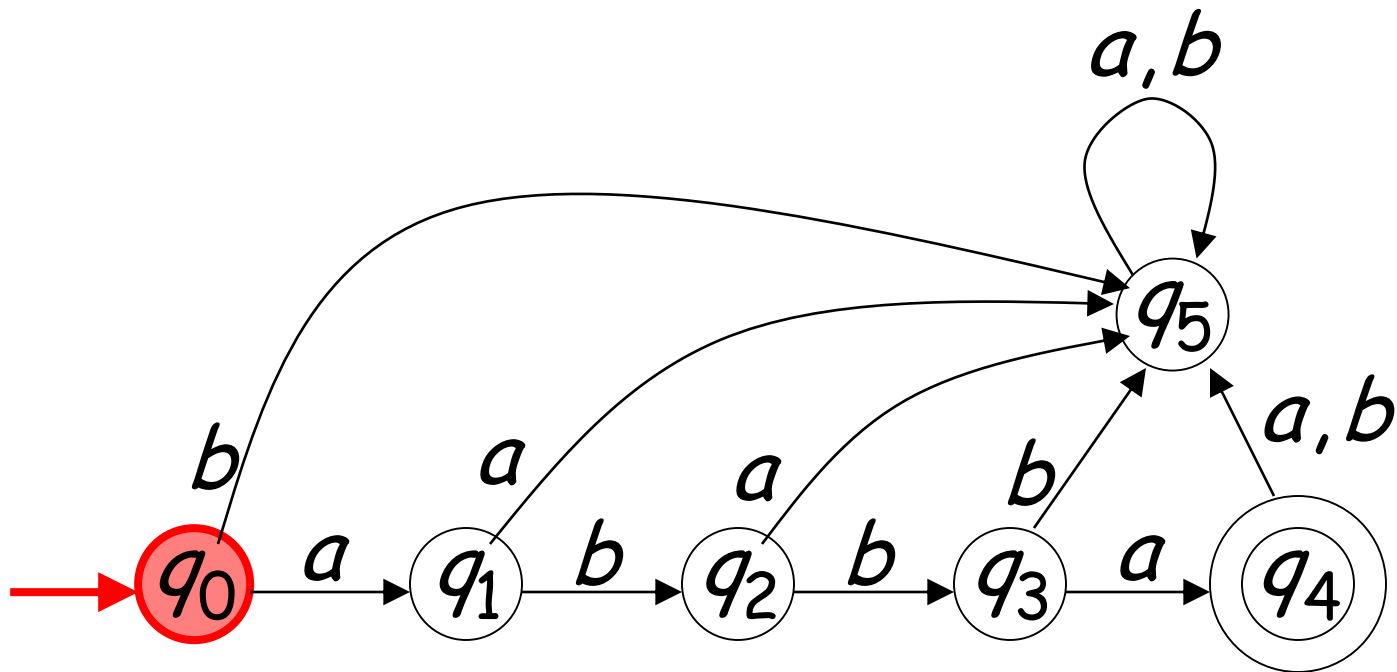
Example

$$\Sigma = \{a, b\}$$



Initial State q_0

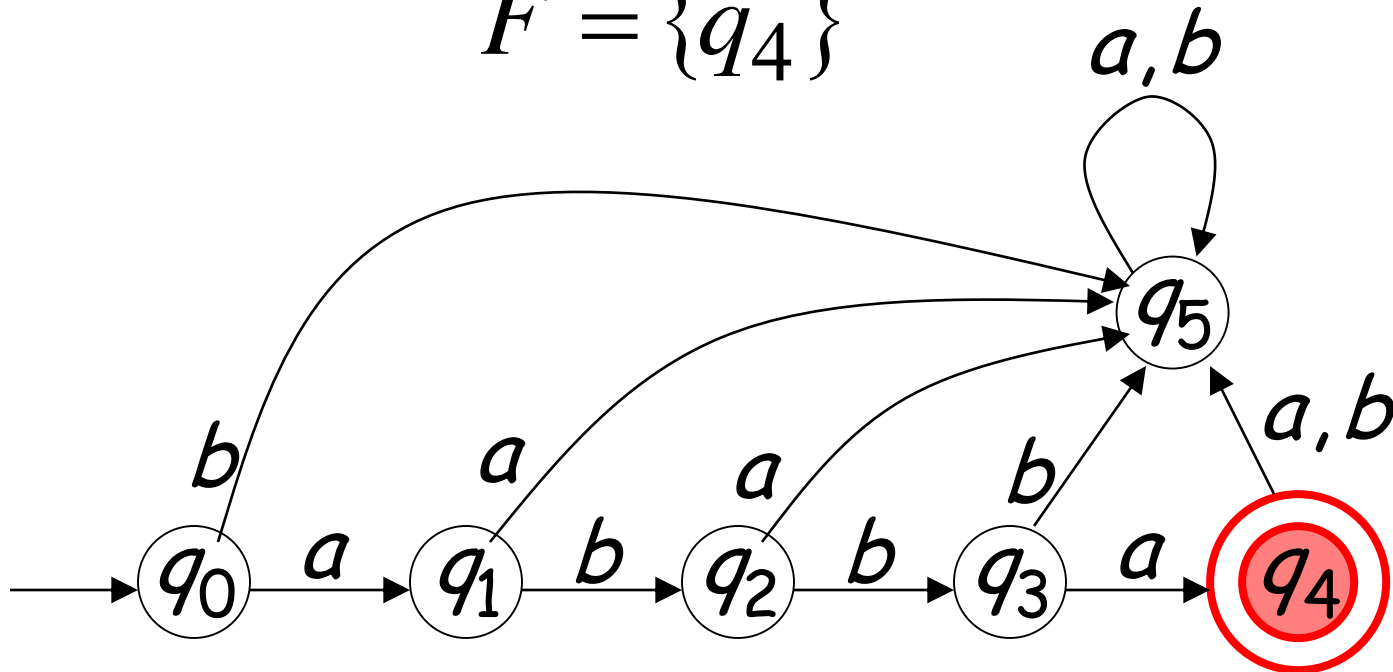
Example



Set of Accepting States $F \subseteq Q$

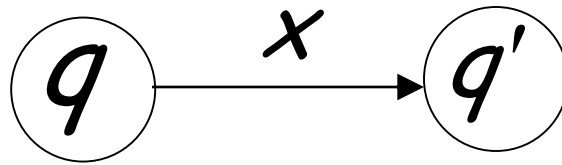
Example

$$F = \{q_4\}$$



Transition Function $\delta : Q \times \Sigma \rightarrow Q$

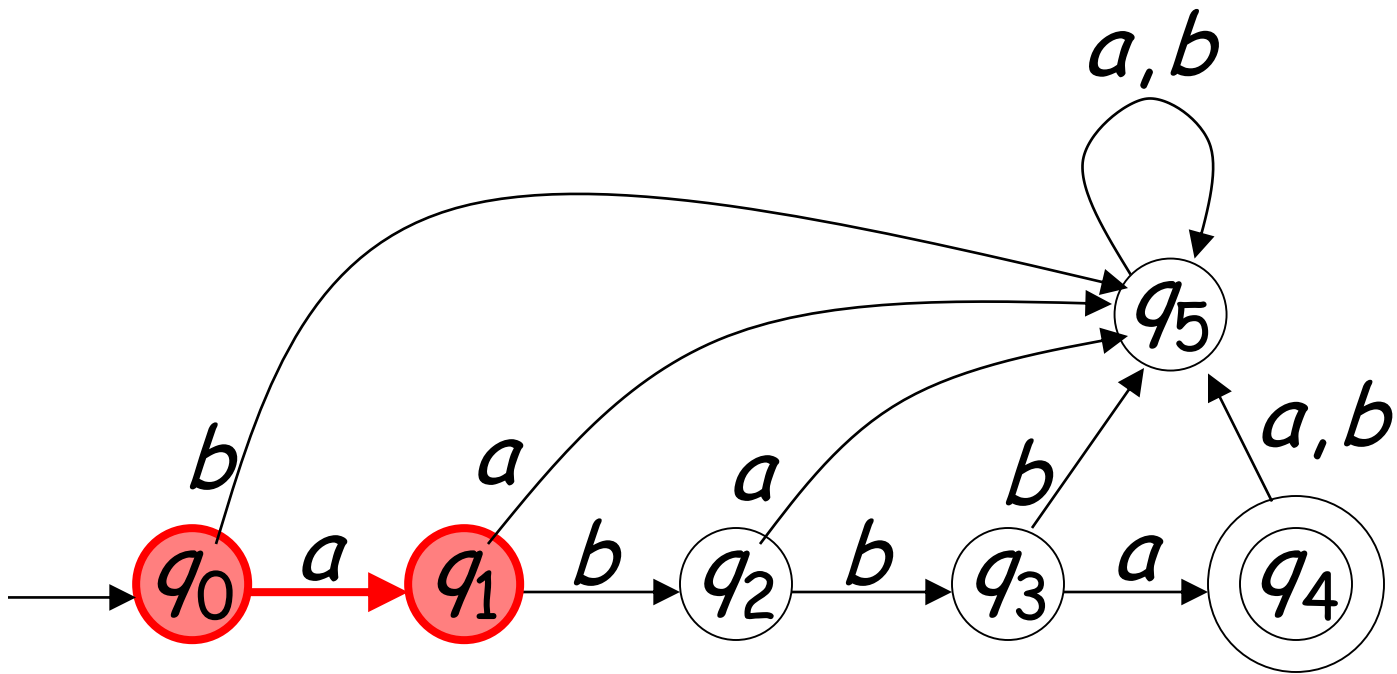
$$\delta(q, x) = q'$$



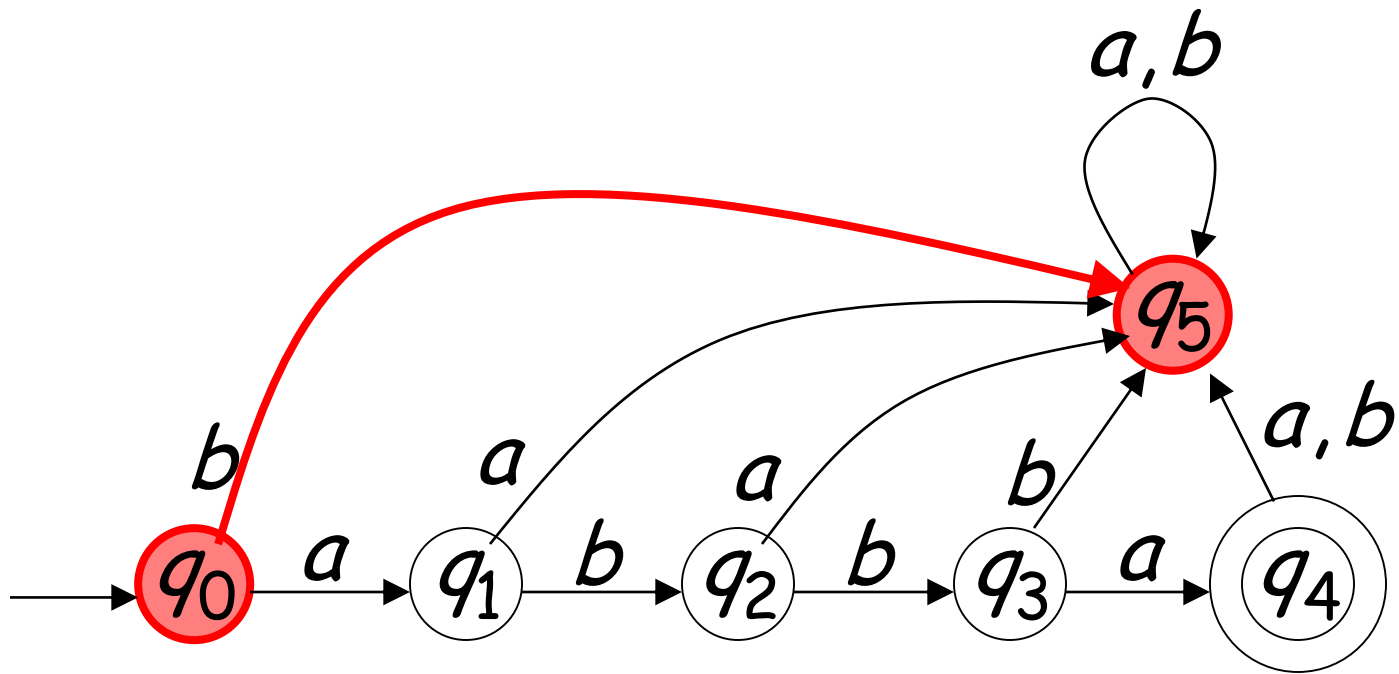
Describes the result of a transition
from state q with symbol x

Example:

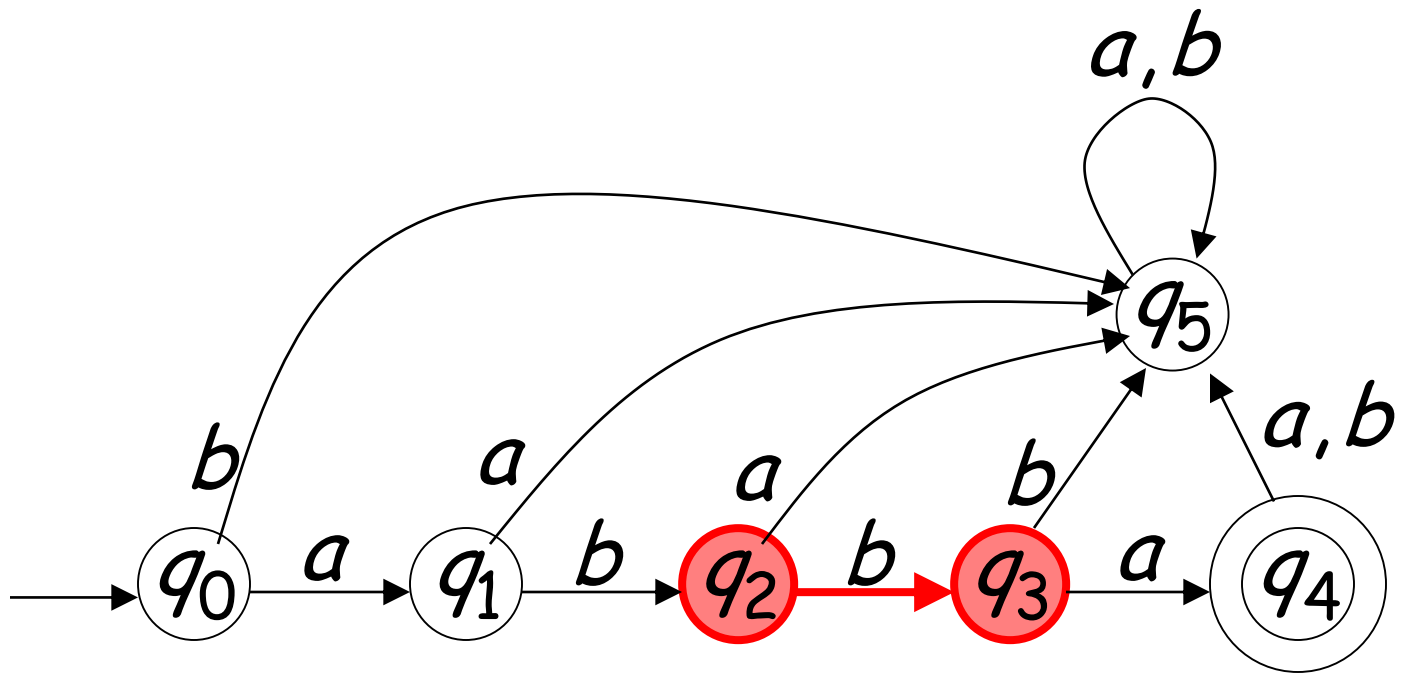
$$\delta(q_0, a) = q_1$$



$$\delta(q_0, b) = q_5$$



$$\delta(q_2, b) = q_3$$

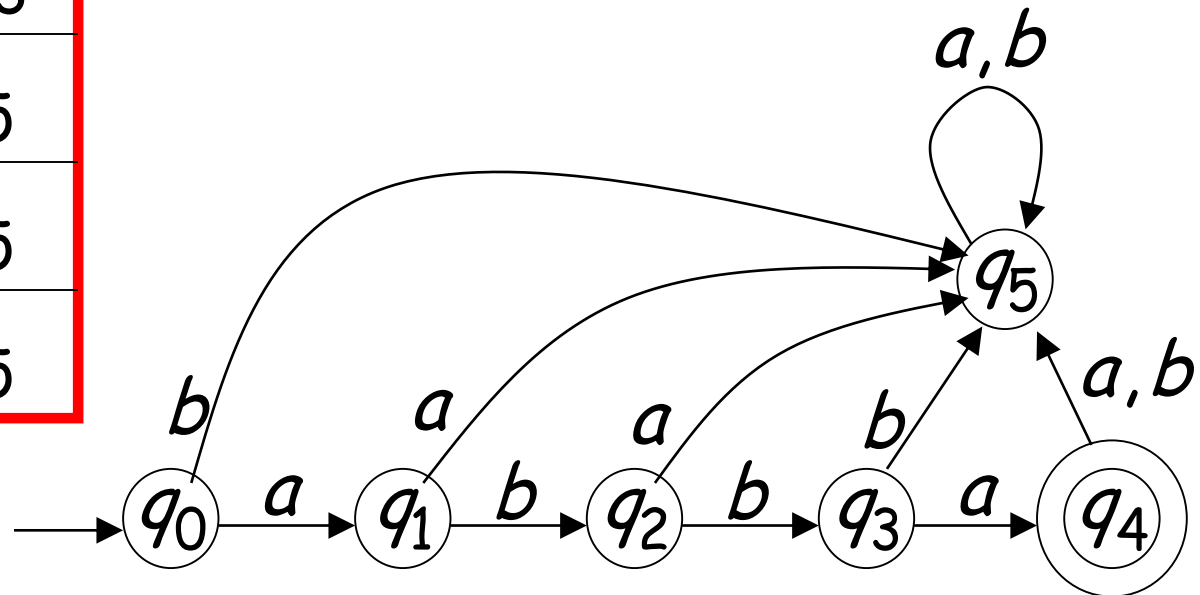


Transition Table for δ

symbols

δ	a	b
q_0	q_1	q_5
q_1	q_5	q_2
q_2	q_5	q_3
q_3	q_4	q_5
q_4	q_5	q_5
q_5	q_5	q_5

states



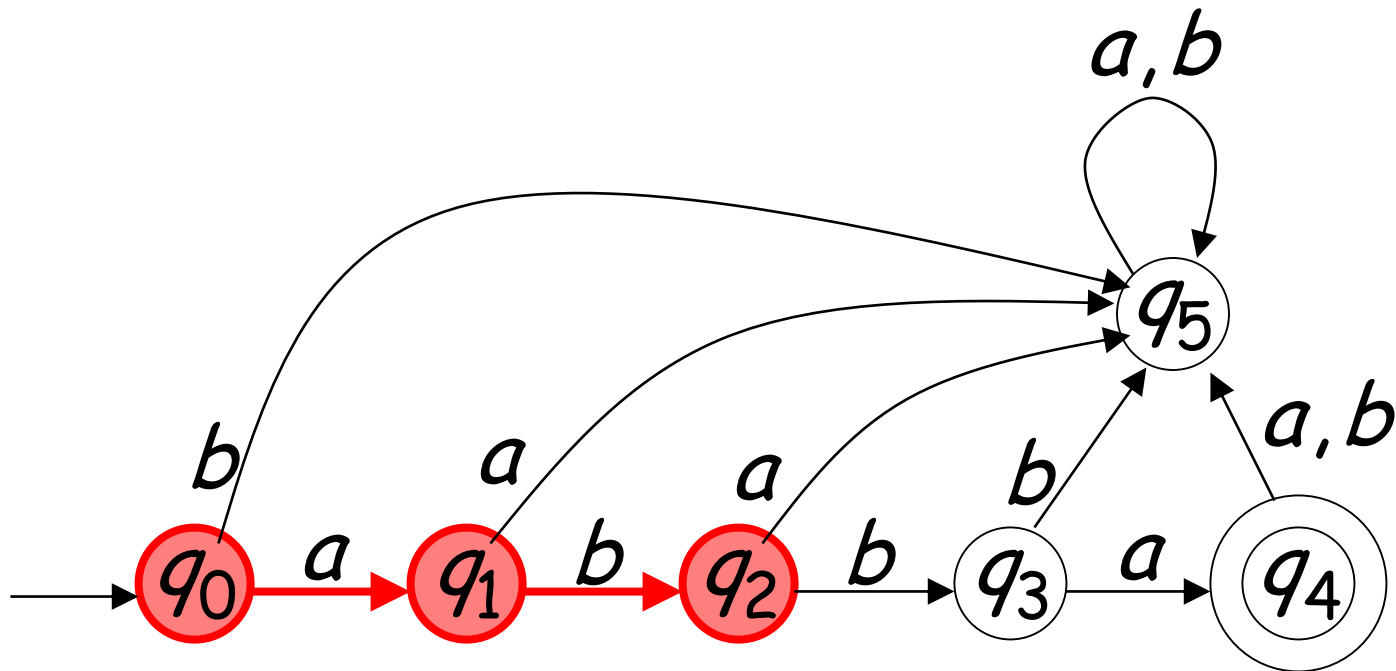
Extended Transition Function

$$\delta^* : Q \times \Sigma^* \rightarrow Q$$

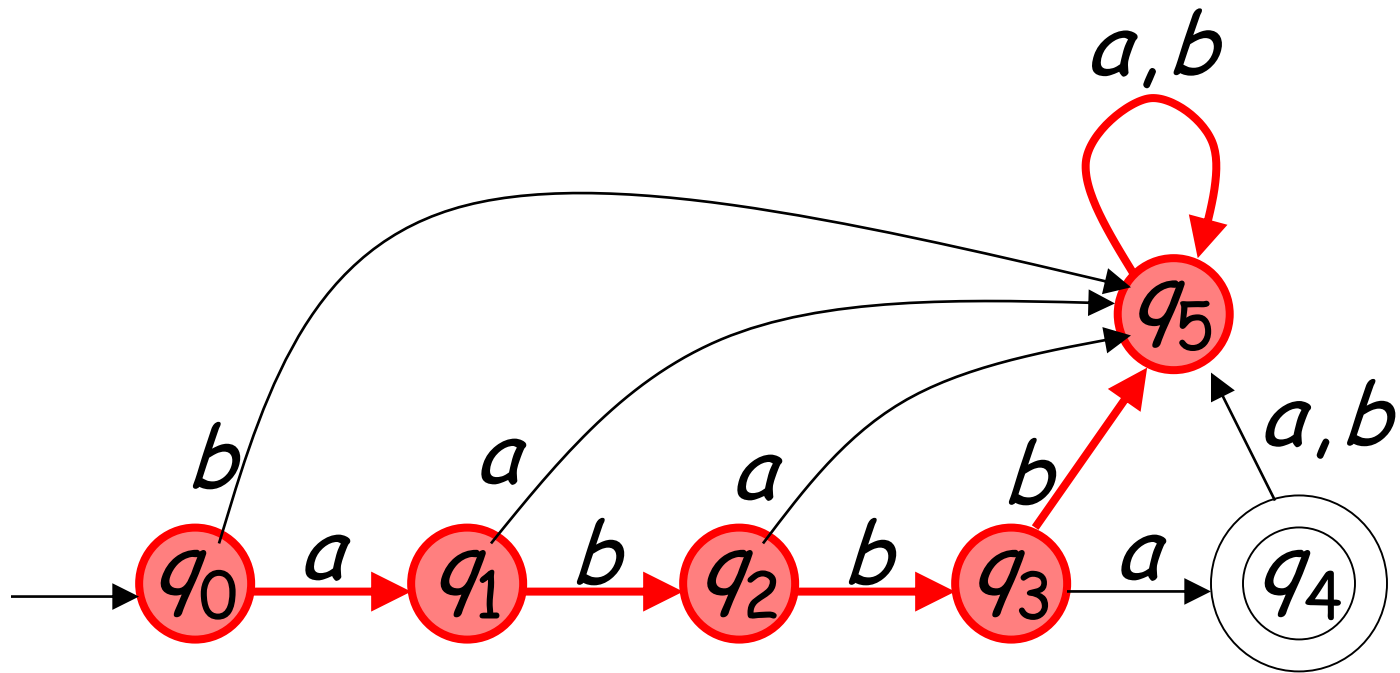
$$\delta^*(q, w) = q'$$

Describes the resulting state
after scanning string w from state q

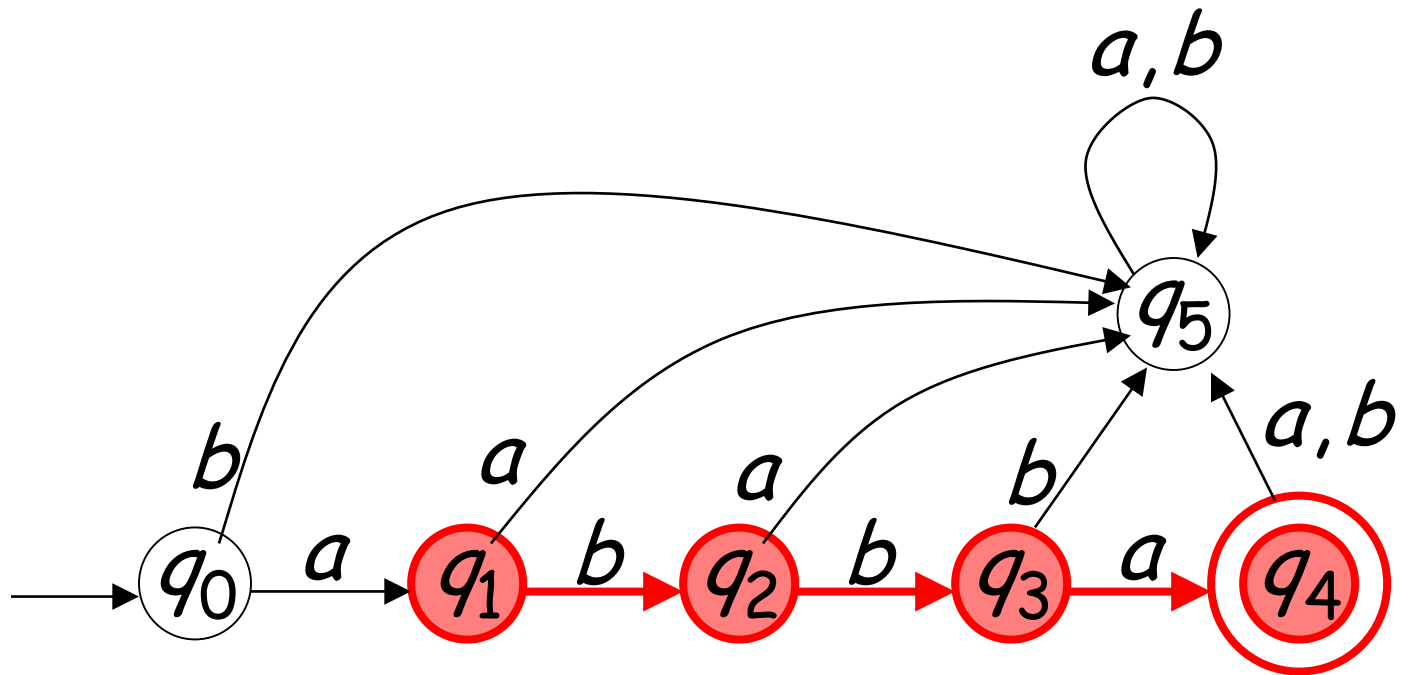
Example: $\delta^*(q_0, ab) = q_2$



$$\delta^*(q_0, abbbaa) = q_5$$



$$\delta^*(q_1, bba) = q_4$$



Special case:

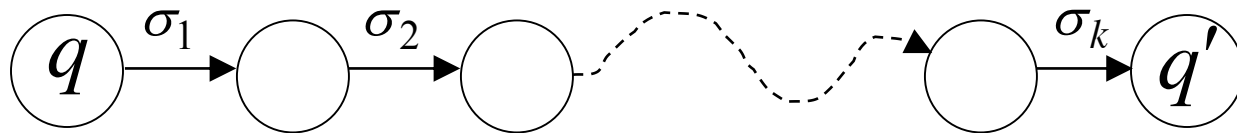
for any state q

$$\delta^*(q, \lambda) = q$$

In general: $\delta^*(q, w) = q'$

implies that there is a walk of transitions

$$w = \sigma_1 \sigma_2 \cdots \sigma_k$$



states may be repeated



Language Accepted by DFA

Language accepted by DFA M :

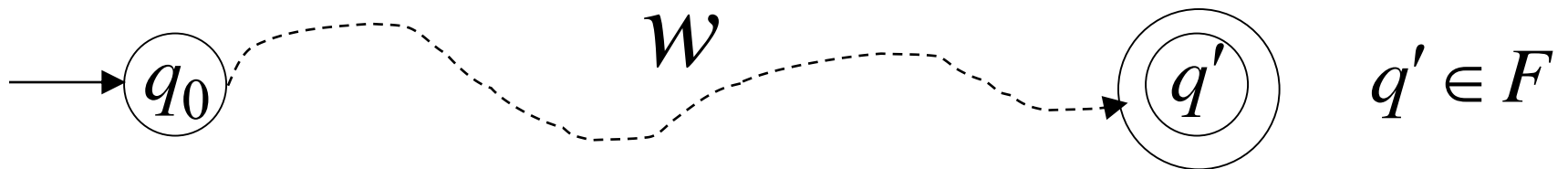
it is denoted as $L(M)$ and contains
all the strings accepted by M

We also say that M recognizes $L(M)$

For a DFA $M = (Q, \Sigma, \delta, q_0, F)$

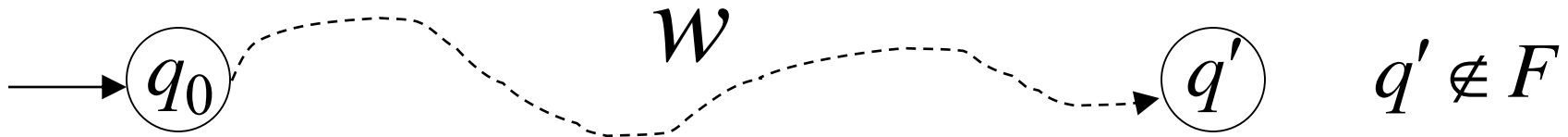
Language accepted by M :

$$L(M) = \{w \in \Sigma^* : \delta^*(q_0, w) \in F\}$$



Language rejected by M :

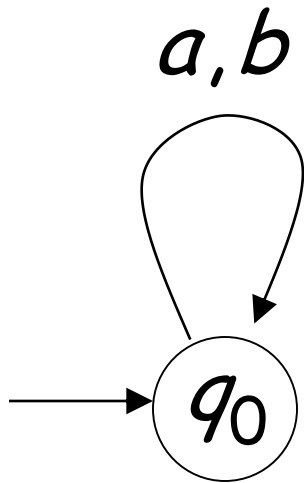
$$\overline{L(M)} = \{w \in \Sigma^* : \delta^*(q_0, w) \notin F\}$$



Notice that we are talking about the complement, as we know, the complement is totally the opposite of the accepting state

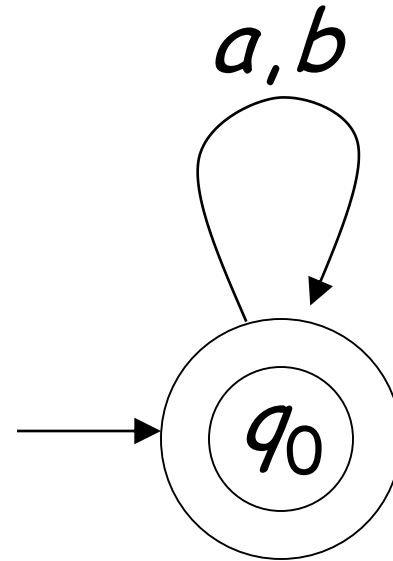
More DFA Examples

$$\Sigma = \{a, b\}$$



$$L(M) = \{ \}$$

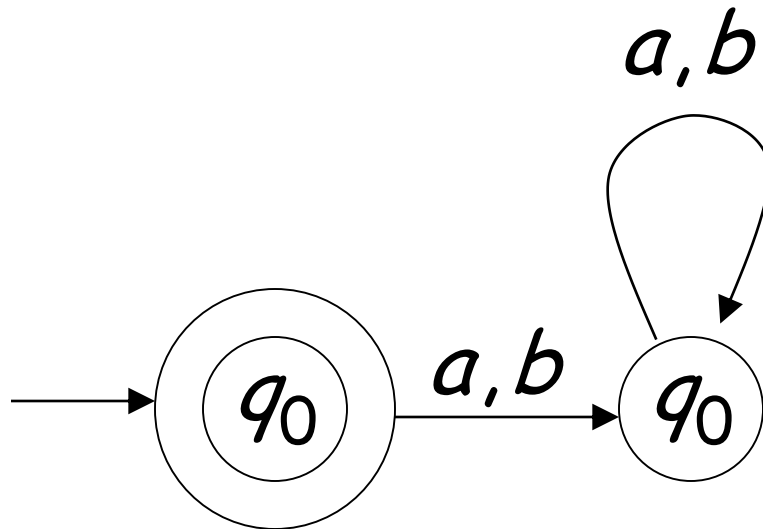
Empty language



$$L(M) = \Sigma^*$$

All strings

$$\Sigma = \{a, b\}$$

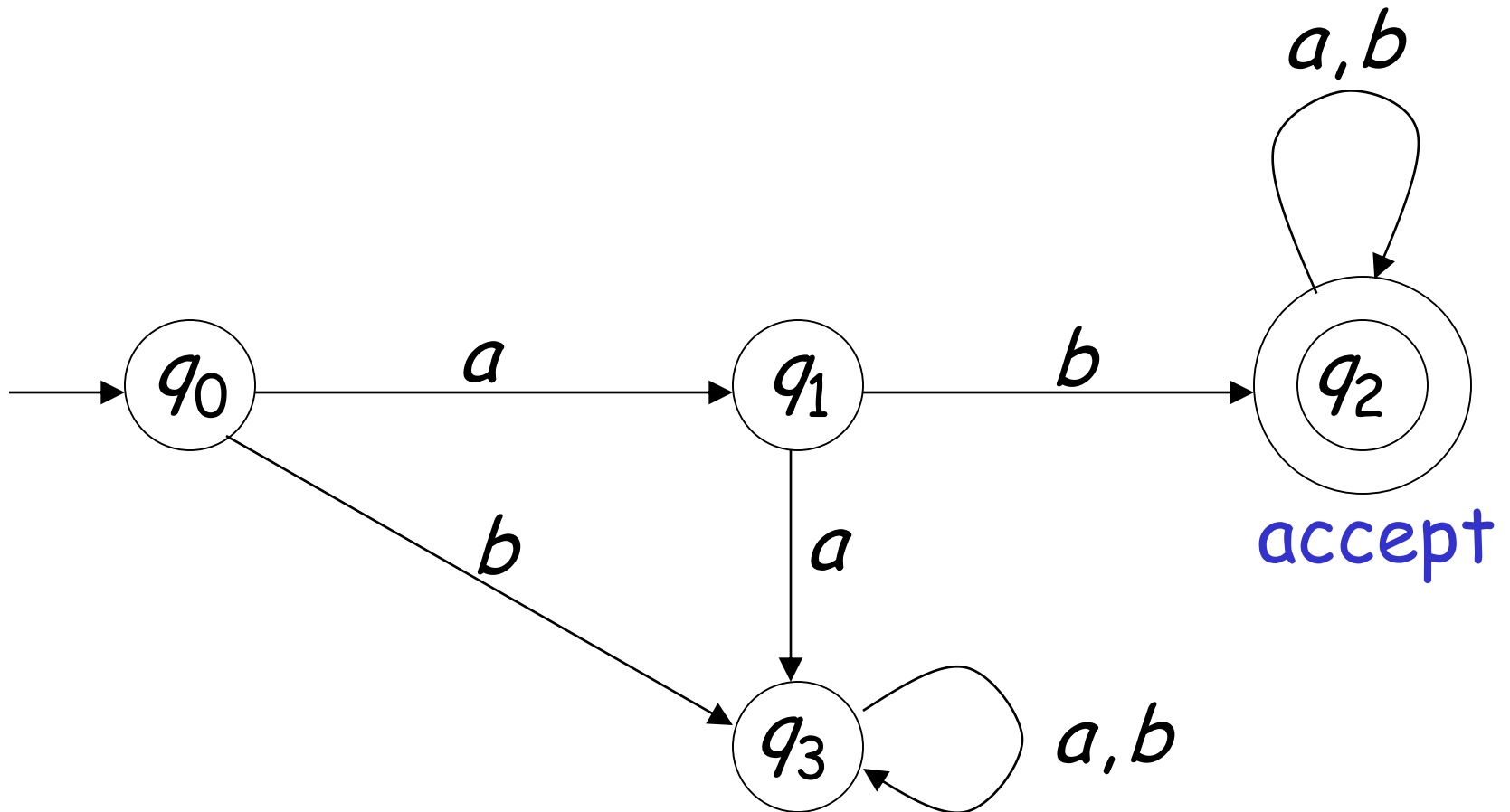


$$L(M) = \{\lambda\}$$

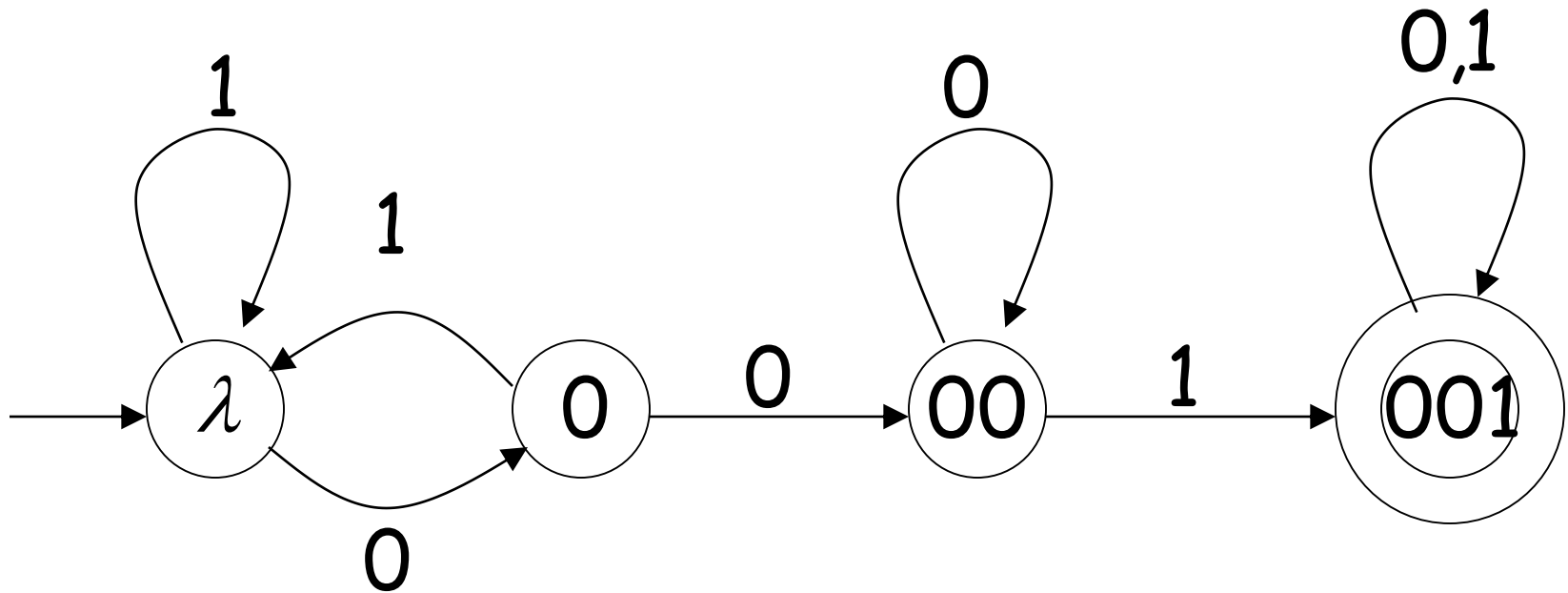
Language of the empty string

$$\Sigma = \{a, b\}$$

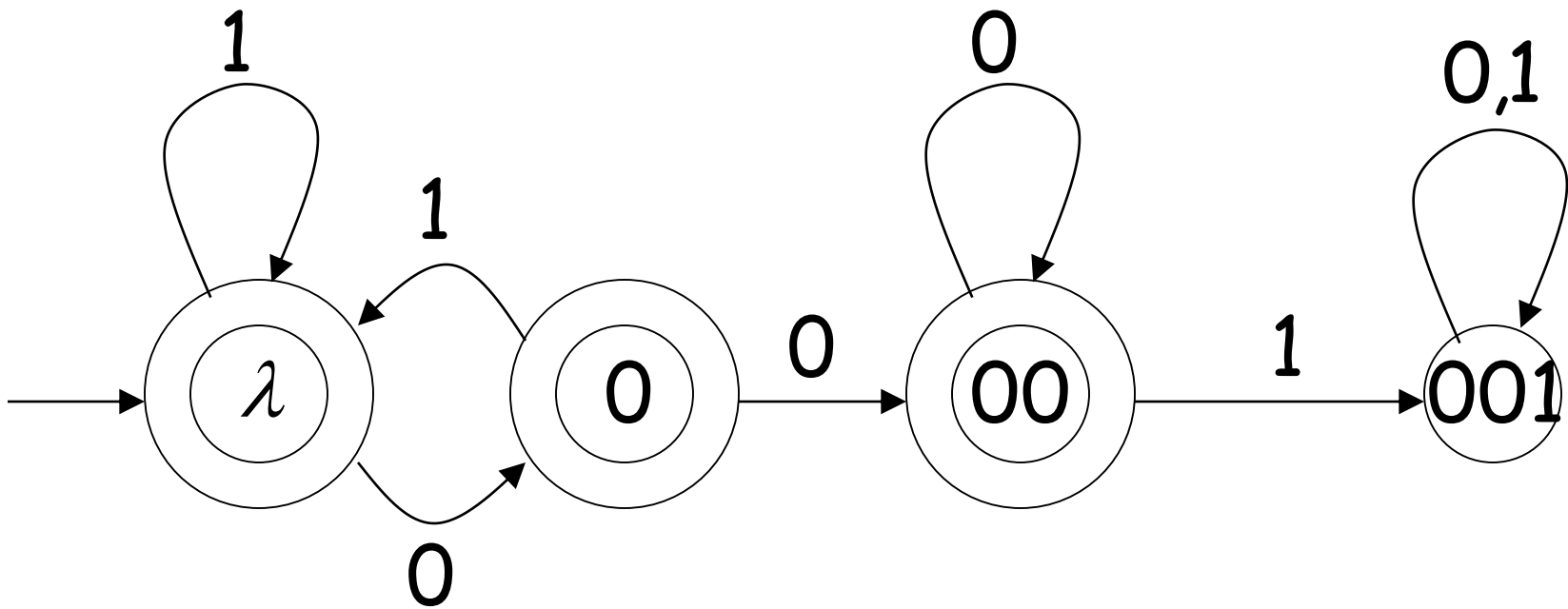
$L(M) = \{ \text{all strings with prefix } ab \}$



$L(M) = \{ \text{all binary strings containing} \\ \text{substring } 001 \}$

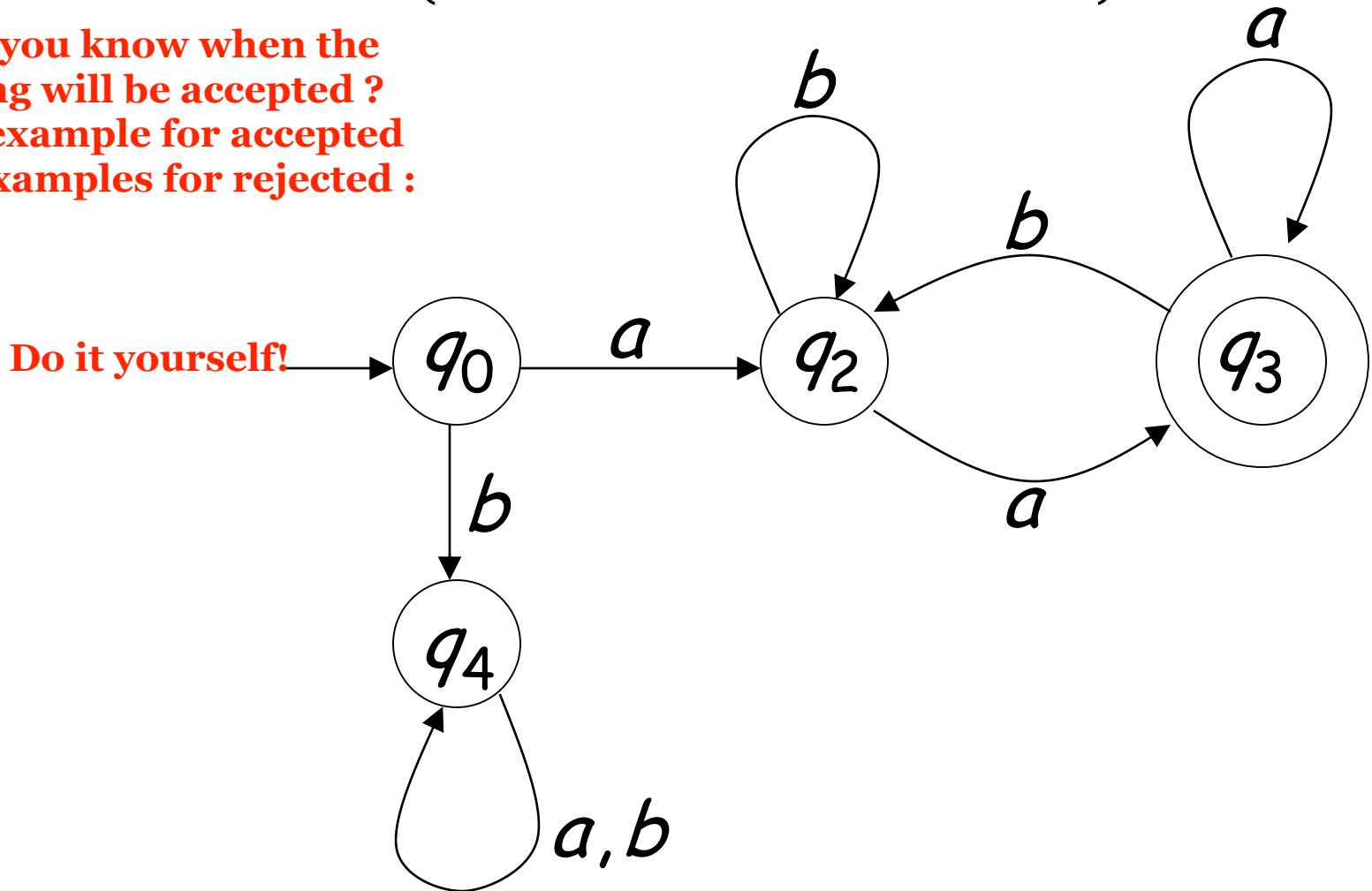


$L(M) = \{ \text{all binary strings without substring } 001 \}$



$$L(M) = \{awa : w \in \{a,b\}^*\}$$

can you know when the string will be accepted?
give example for accepted
and examples for rejected:



Regular Languages

Definition:

A language L is **regular** if there is a DFA M that accepts it ($L(M) = L$)

The languages accepted by all DFAs form the family of **regular languages**

Example regular languages:

$\{abba\}$ $\{\lambda, ab, abba\}$

$\{a^n b : n \geq 0\}$ $\{awa : w \in \{a, b\}^*\}$

$\{\text{all strings in } \{a, b\}^* \text{ with prefix } ab\}$

$\{\text{all binary strings without substring } 001\}$

$\{x : x \in \{1\}^* \text{ and } x \text{ is even}\}$

$\{\}$ $\{\lambda\}$ $\{a, b\}^*$

There exist DFAs that accept these languages (see previous slides).

There exist languages which are not Regular:

$$L = \{a^n b^n : n \geq 0\}$$

$$\text{ADDITION} = \{x + y = z : x = 1^n, y = 1^m, z = 1^k, \\ n + m = k\}$$

There are no DFAs that accept these languages

(we will prove this in a later class)

Think about this until next class !